

Formal Techniques – 2018-06-28

Exercise 1. Informally describe the four CTL formulae $AF\phi$, $AG\phi$, $EF\phi$, $EG\phi$ (where ϕ is atomic), providing for each one a brief description (1-3 lines), and one example where it holds.

Exercise 2. Consider the following protocol excerpt written in the applied-pi notation.

(in W . ! . out $f(W)$. () | in X . out $h(X)$. in Y . out $g(X, Y)$. ())

Apply the control flow analysis to the protocol above, generating a tree automaton to over-approximate the message flow, as done by function $gen(\dots)$. Provide a list of states for such automaton and the transitions among them. Make each state clearly related to a part of the protocol above.

Exercise 3. Formalize the following cryptographic protocol fragment using the applied-pi notation.

Initially, Alice knows symmetric keys k_1, k_2 . Further, Alice and Bob share a symmetric key k_s .

- 1) Alice sends k_1 to Bob, encrypting it with k_s .
- 2) Then, Bob generates a fresh nonce N , and sends N to Alice, encrypting it with k_1 .
- 3) Alice answers by generating a message containing two items: the key k_2 and the hash of N encrypted with k_2 . The whole message (a pair) is encrypted using k_s .
- 4) Bob finally retrieves the hash of N from the received message, and sends such hash to Alice (without encrypting it).

Exercise 4. Formally prove the following formula exploiting the Curry-Howard isomorphism.

$$\forall p, q, r, s : \text{Prop. } ((p \wedge q) \rightarrow s) \rightarrow ((q \vee r) \rightarrow (r \vee (p \rightarrow s)))$$

Exercise 5. Let $\alpha : \mathcal{C} \xleftarrow{\quad} \mathcal{A} : \gamma$ be a Galois connection between two CLs \mathcal{C}, \mathcal{A} .

1. Prove that $\gamma \circ \alpha \circ \gamma = \gamma$.
2. Define $\delta : \mathcal{A} \rightarrow \mathcal{A}$ as the function $\delta(a) = \prod \{a' \mid \gamma(a') = \gamma(a)\}$. Prove that $\delta = \alpha \circ \gamma$.

Exercise 6.

1. Below, we write $R \subseteq^{rs} X^2$ when R is a binary relation over set X , and R is reflexive ($\forall x \in X. xRx$) and symmetric ($\forall x, y \in X. xRy \implies yRx$).

For any $R \subseteq^{rs} X^2$, define $X_R = \{A \subseteq X \mid \forall x, y \in A. xRy\}$.

Prove that (X_R, \subseteq) is a DCPO with $\sqcup = \bigcup$ and $\perp = \emptyset$.

2. Let $A \sim_R B$ mean $\forall a \in A, b \in B. aRb$. For any two relations $R, S \subseteq^{rs} X^2$, define

$$St(R, S) = \left\{ f : X_R \rightarrow X_S \mid \begin{array}{l} f \text{ Scott-continuous } \wedge \\ \forall A, B \in X_R. A \sim_R B \implies f(A \cap B) = f(A) \cap f(B) \end{array} \right\}$$

Prove that, if $R, S, T \subseteq^{rs} X$, $f \in St(R, S)$, and $g \in St(S, T)$, then $g \circ f \in St(R, T)$.