## Formal Techniques – 2023-06-23

**Exercise 1.** Informally describe the four CTL formulae  $AF\phi$ ,  $AG\phi$ ,  $EF\phi$ ,  $EG\phi$  (where  $\phi$  is atomic), providing for each one a brief description (1-3 lines), and one example where it holds.

**Exercise 2.** Consider the following protocol excerpt written in the applied-pi notation.

! . new X . out f(X) . (in Y . out g(X, Y) . () | in Z . in W . out h(W, Z) . ())

Apply the control flow analysis to the protocol above, generating a tree automaton to over-approximate the message flow, as done by function gen(...). Provide a list of states for such automaton and the transitions among them. Make each state clearly related to a part of the protocol above.

**Exercise 3.** Consider the following tree automaton

and the rewriting rule

 $\mathsf{dec}(\mathsf{enc}(M,K),K) \Rightarrow M$ 

Apply the completion algorithm to the above automaton, building an overapproximation for the languages associated to its states which is closed under rewriting. Assuming @a models the set of messages being exchanged over a public channel, state what can be concluded about the secrecy of m.

**Exercise 4.** Formally prove the following formula exploiting the Curry-Howard isomorphism.

$$\forall p, q, r, s : \mathsf{Prop.} \ (p \lor q) \to [(p \to (q \to s)) \to [(p \to s) \lor (q \to s))]$$

**Exercise 5.** Let  $\mathcal{A}$  be a CL, and let  $f, g, h : \mathcal{A} \to \mathcal{A}$  be three Scott-continuous functions, with  $h(a) = f(a) \sqcap g(a)$  for all  $a \in \mathcal{A}$ . Writing fix for the least fixed point operator, let

$$x = \operatorname{fix}(h)$$
  $y = \bigsqcup_{n \ge 0} f^n(x)$ 

- 1. [10%] Show that  $\bigsqcup_{n>0} f^n(x)$  is the supremum of a directed set.
- 2. [90%] Show that y = fix(f).
- (Bonus) Show that, in general, when b is an arbitrary element of A, then ⊔<sub>n>0</sub> f<sup>n</sup>(b) does not always have to be a fixed point of f.