Computability Assignment Year 2012/13 - Number 1

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1 Question

Define a binary property p(x, y) over natural numbers such that we have both

- 1. $\forall x \in \mathbb{N} . \exists y \in \mathbb{N} . p(x, y)$
- 2. $\neg \exists y \in \mathbb{N} . \forall x \in \mathbb{N} . p(x, y)$

Provide a definition for p, and a proof for the above claims.

1.1 Answer

Define p(x, y) as follow p(x, y) = (x = y) in this case both 1. and 2. are satisfied.

1.1.1 $\forall x \in \mathbb{N} . \exists y \in \mathbb{N} . p(x, y)$

RZ: ok, but you don't really need induction to prove x = x

The property is satisfied for x = 0 because exists an y that makes it true y = 0 in fact p(0,0) = (0 = 0) that it's true.

Now we assume that the theorem it's true for the generic x so we see that y = x makes the preposition true.

The last step show us that for x + 1 the the preposition is true in fact is satisfied by y = x + 1:

if y = x + 1

p(y, x + 1) = (y = x + 1) we sobstitute y by the rules above

p(x+1, x+1) = (x+1 = x+1) so is satisfied.

Inductively we have proof the first claim.

1.1.2 $\neg \exists y \in \mathbb{N}. \forall x \in \mathbb{N}. p(x, y)$

If we choose y = 0 and x = 0 the property is verified but not for y = 0 and x = 1 so it isn't verified for all the x.

We can now take the generic y, the property is satisfied for x by the rules x = y but not for x + 1 by the same rules in fact if x = y, $x + 1 \neq y$. For y + 1 the rules is the same as above in fact if y = x: y + 1 = x + 1 but $y + 1 \neq x + 2$ so $\neg \exists y \in \mathbb{N}. \forall x \in \mathbb{N}. p(x, y)$ is satisfied.