Computability Assignment Year 2012/13 - Number 1

Please keep this file anonymous: do not write your name inside this file. More information about assignments at http://disi.unitn.it/~zunino/teaching/computability/assignments

1 Question

Define a binary property p(x, y) over natural numbers such that we have both

- 1. $\forall x \in \mathbb{N} : \exists y \in \mathbb{N} : p(x, y)$
- 2. $\neg \exists y \in \mathbb{N} . \forall x \in \mathbb{N} . p(x, y)$

Provide a definition for p, and a proof for the above claims.

1.1 Answer

I Start defining the property p as the greater function > between y and x (y > x).

Now I can verify the proposition 1 by induction. First of all I define y = succ(x) where succ(x) = x + 1.

RZ: you don't really need induction here, it's enough to say that p(x, x + 1) is true since x + 1 > x.

Base case

Setting x = 0 we obtain that y = succ(x) = 1. The property p(0, 1) is verified since 1 > 0 is true.

Inductive Step

We consider valid the property till x = n. We have to prove that p is also valid for x = n + 1.

We get that y = succ(x) = n + 1 + 1, and so p(n + 1 + 1, n + 1). Since n + 1 + 1 > n + 1, then the property p is verified for all x. It follows that the first proposition is verified.

Now I have to verify the proposition 2, by induction again.

This proposition can be rewritten as follows: $\neg \exists y \in \mathbb{N} . \forall x \in \mathbb{N} . p(x, y) \equiv \forall y \in \mathbb{N} . \neg \forall x \in \mathbb{N} . p(x, y) \equiv \forall y \in \mathbb{N} . \exists x \in \mathbb{N} . \neg p(x, y)$

Now I set $\neg p(x, y) \equiv q(x, y)$ which can be interpreted as $x \leq y$.

RZ: as above, you don't need induction. It's enough to say that if you pick x=0, you have q(x, y) for all y since $0 \le y$.

Base case

Setting y = 0 and x = 0 the property q(0,0) is verified since $0 \le 0$ is true.

Inductive Step

We consider valid the property till y = n. We have to prove that q is also valid for y = n + 1.

We set x = 0, so we obtain q(0, n + 1). Since $0 \le n + 1$, then the property q is verified for all y.

Since the property q is always valid we can infer, moving in a reverse manner, that the property p is always true.

It follows that the second proposition is verified.