Fundamentals of Artificial Intelligence Chapter 10: Classical Planning

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- Basics on Planning
 - The Problem
 - The PDDL Language
- Search Strategies and Heuristics
 - Forward and Backward Search
 - Heuristics
- Planning Graphs, Heuristics and Graphplan
 - Planning Graphs
 - Heuristics Driven by Planning Graphs
 - The Graphplan Algorithm
- Other Approaches (hints)
 - Planning as SAT Solving
 - Planning as FOL Inference



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Automated Planning

- Planning is both:
 - an application per se
 - a common activity in many applications
 (e.g. design & manufacturing, scheduling, robotics,...)
- Similar to problem-solving agents (Ch.03), but with factored/structured representation of states
- "Classical" Planning (this chapter): fully observable, deterministic, static environments with single agents

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Automated Planning [cont.]

Automated Planning

- Given:
 - an initial state
 - a set of actions you can perform
 - a (set of) state(s) to achieve (goal)
- Find:
 - a plan: a partially- or totally-ordered set of actions needed to achieve the goal from the initial state

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Decidability and Complexity

- PlanSAT: the question of whether there exists any plan that solves a planning problem
 - decidable for classical planning
 - with function symbols, the number of states becomes infinite
 - ⇒ undecidable
 - in PSPACE
- Bounded PlanSAT: the question of whether there exists any plan that of a given length k or less
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- A state is a conjunction of fluents: ground, function-less atoms
 - ex: Poor ∧ Unknown, At(Truck₁, Melbourne) ∧ At(Truck₂, Sydney)
 - ex of non-fluents: At(x, y) (non ground), $\neg Poor$ (negated), At(Father(Fred), Sydney) (not function-less)
 - closed-world assumption: all non-mentioned fluents are false
 - unique names assumption: distinct names refer to distinct objects
- Actions are described by a set of action schemata
 - concise description: describe which fluent change
 - ⇒ the other fluents implicitly maintain their values
- Action Schema: consists in action name, a list of variables in the schema, the precondition, the effect (aka postcondition)
 - precondition and effect are conjunctions of literals (positive or negated atomic sentences)
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PDDL: Example

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• Action schema:
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Action(Fly(p, from, to), PRECOND : At(p, from) \land Plane(p) \land Airport(from) \land Airport(to)

EFFECT : \neg At(p, from) \land At(p, to))
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• Action instantiation:

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Action(Fly(P_1, SFO, JFK)),

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- Precondition: must hold to ensure the action can be executed
 - defines the states in which the action can be executed
 - action is applicable in state s if the preconditions are satisfied by s
- Effect: represent the effects of the action on the world
 - defines the result of executing the action
- Add list (ADD(a)): (the fluents in) the positive literals in the action's effects
 - ex: $\{At(p, to)\}$
- Delete list (DEL(a)): (the fluents in) the negative literals in the action's effects
 - ex: {*At*(*p*, *from*)}
- Result of action a in state s: RESULT(s,a) $\stackrel{\text{def}}{=}$ (s\DEL(a) \cup ADD(a))
 - start from s
 - remove the fluents that appear as negative literals in effect
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 - ex: $Fly(P_1, SFO, JFK) \Longrightarrow \text{remove } At(P_1, SFO), \text{ add } At(P_1, JFK)$

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• $s : At(P_1, SFO) \land Plane(P_1) \land Airport(SFO) \land Airport(JFK) \land ...$

 $\implies s': At(P_1, JFK) \land Plane(P_1) \land Airport(SFO) \land Airport(JFK) \land ...$

Sometimes we want to propositionalize a PDDL problem: replace each action schema with a set of ground actions.

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Ex: ...At_P₁_SFO ∧ Plane_P₁ ∧ Airport_SFO ∧ Airport_JFK)...

Time in PDDL

- Fluents do not explicitly refer to time
- Times and states are implicit in the action schemata:
 - the precondition always refers to time t
 - the effect to time t+1.

- A set of action schemata defines a planning domain
- PDDL problem: a planning domain, an initial state and a goal
 - the initial state is a conjunction of ground atoms (positive literals)
 - closed-world assumption: any not-mentioned atoms are false
 - the goal is a conjunction of literals (positive or negative
 - may contain variables, which are implicitly existentially quantified
 - a goal g may represent a set of states (the set of states entailing g)
- \bullet Ex: goal: $At(p, SFO) \land Plane(p)$:
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Planning as a search problem

All components of a search problem

- an initial state
- an ACTIONS function
- a RESULT function
- and a goal test

Example: Air Cargo Transport

```
Init(At(C_1, SFO) \wedge At(C_2, JFK) \wedge At(P_1, SFO) \wedge At(P_2, JFK)
    \wedge Cargo(C_1) \wedge Cargo(C_2) \wedge Plane(P_1) \wedge Plane(P_2)
    \land Airport(JFK) \land Airport(SFO)
Goal(At(C_1, JFK) \wedge At(C_2, SFO))
Action(Load(c, p, a),
  PRECOND: At(c, a) \wedge At(p, a) \wedge Cargo(c) \wedge Plane(p) \wedge Airport(a)
  EFFECT: \neg At(c, a) \land In(c, p)
Action(Unload(c, p, a),
  PRECOND: In(c, p) \wedge At(p, a) \wedge Cargo(c) \wedge Plane(p) \wedge Airport(a)
  EFFECT: At(c, a) \land \neg In(c, p)
Action(Fly(p, from, to),
  PRECOND: At(p, from) \wedge Plane(p) \wedge Airport(from) \wedge Airport(to)
  EFFECT: \neg At(p, from) \land At(p, to)
```

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One solution: $[Load(C_1, P_1, SFO), Fly(P_1, SFO, JFK), Unload(C_1, P_1, JFK), Load(C_2, P_2, JFK), Fly(P_2, JFK, SFO), Unload(C_2, P_2, SFO)]$

Example: Air Cargo Transport

```
Init(At(C_1, SFO) \wedge At(C_2, JFK) \wedge At(P_1, SFO) \wedge At(P_2, JFK)
    \wedge Cargo(C_1) \wedge Cargo(C_2) \wedge Plane(P_1) \wedge Plane(P_2)
    \land Airport(JFK) \land Airport(SFO)
Goal(At(C_1, JFK) \wedge At(C_2, SFO))
Action(Load(c, p, a),
  PRECOND: At(c, a) \wedge At(p, a) \wedge Cargo(c) \wedge Plane(p) \wedge Airport(a)
  EFFECT: \neg At(c, a) \land In(c, p)
Action(Unload(c, p, a),
  PRECOND: In(c, p) \wedge At(p, a) \wedge Cargo(c) \wedge Plane(p) \wedge Airport(a)
  EFFECT: At(c, a) \land \neg In(c, p)
Action(Fly(p, from, to),
  PRECOND: At(p, from) \wedge Plane(p) \wedge Airport(from) \wedge Airport(to)
  EFFECT: \neg At(p, from) \land At(p, to)
```

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One solution: $[Load(C_1, P_1, SFO), Fly(P_1, SFO, JFK), Unload(C_1, P_1, JFK), Load(C_2, P_2, JFK), Fly(P_2, JFK, SFO), Unload(C_2, P_2, SFO)]$

Example: Spare Tire Problem

```
Init(Tire(Flat) \land Tire(Spare) \land At(Flat, Axle) \land At(Spare, Trunk))
Goal(At(Spare, Axle))
Action(Remove(obj, loc),
  PRECOND: At(obj, loc)
  EFFECT: \neg At(obj, loc) \land At(obj, Ground)
Action(PutOn(t, Axle),
   PRECOND: Tire(t) \land At(t, Ground) \land \neg At(Flat, Axle)
   EFFECT: \neg At(t, Ground) \land At(t, Axle)
Action(LeaveOvernight,
   PRECOND:
   EFFECT: \neg At(Spare, Ground) \land \neg At(Spare, Axle) \land \neg At(Spare, Trunk)
            \wedge \neg At(Flat, Ground) \wedge \neg At(Flat, Axle) \wedge \neg At(Flat, Trunk)
```

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(We assume that the car is parked in a particularly bad neighborhood, so that the effect of leaving it overnight is that the tires disappear.)

One solution: [Remove(Flat, Axle), Remove(Spare, Trunk), PutOn(Spare, Axle)

Example: Spare Tire Problem

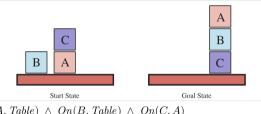
```
Init(Tire(Flat) \land Tire(Spare) \land At(Flat, Axle) \land At(Spare, Trunk))
Goal(At(Spare, Axle))
Action(Remove(obj, loc),
  PRECOND: At(obj, loc)
  EFFECT: \neg At(obj, loc) \land At(obj, Ground)
Action(PutOn(t, Axle),
   PRECOND: Tire(t) \land At(t, Ground) \land \neg At(Flat, Axle)
   EFFECT: \neg At(t, Ground) \land At(t, Axle)
Action(LeaveOvernight,
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   EFFECT: \neg At(Spare, Ground) \land \neg At(Spare, Axle) \land \neg At(Spare, Trunk)
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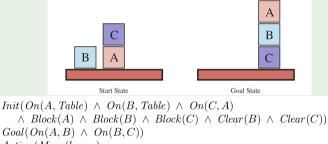
Example: Blocks World



$$\begin{split} &\operatorname{Init}(\operatorname{On}(A,\operatorname{Table})\,\wedge\,\operatorname{On}(B,\operatorname{Table})\,\wedge\,\operatorname{On}(C,A) \\ &\wedge\,\operatorname{Block}(A)\,\wedge\,\operatorname{Block}(B)\,\wedge\,\operatorname{Block}(C)\,\wedge\,\operatorname{Clear}(B)\,\wedge\,\operatorname{Clear}(C)) \\ &\operatorname{Goal}(\operatorname{On}(A,B)\,\wedge\,\operatorname{On}(B,C)) \\ &\operatorname{Action}(\operatorname{Move}(b,x,y), \\ &\operatorname{PRECOND:}\,\operatorname{On}(b,x)\,\wedge\,\operatorname{Clear}(b)\,\wedge\,\operatorname{Clear}(y)\,\wedge\,\operatorname{Block}(b)\,\wedge\,\operatorname{Block}(y)\,\wedge\,\\ &(b\neq x)\,\wedge\,(b\neq y)\,\wedge\,(x\neq y), \\ &\operatorname{Effect:}\,\operatorname{On}(b,y)\,\wedge\,\operatorname{Clear}(x)\,\wedge\,\neg\operatorname{On}(b,x)\,\wedge\,\neg\operatorname{Clear}(y)) \\ &\operatorname{Action}(\operatorname{MoveToTable}(b,x), \\ &\operatorname{PRECOND:}\,\operatorname{On}(b,x)\,\wedge\,\operatorname{Clear}(b)\,\wedge\,\operatorname{Block}(b)\,\wedge\,(b\neq x), \\ &\operatorname{Effect:}\,\operatorname{On}(b,\operatorname{Table})\,\wedge\,\operatorname{Clear}(x)\,\wedge\,\neg\operatorname{On}(b,x)) \end{split}$$

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Example: Blocks World



 $\begin{array}{l} Action(Move(b,x,y), \\ \text{PRECOND: } On(b,x) \ \land \ Clear(b) \ \land \ Clear(y) \ \land \ Block(b) \ \land \ Block(y) \ \land \\ (b \neq x) \ \land \ (b \neq y) \ \land \ (x \neq y), \\ \text{EFFECT: } On(b,y) \ \land \ Clear(x) \ \land \ \neg On(b,x) \ \land \ \neg Clear(y)) \\ Action(MoveToTable(b,x), \\ \text{PRECOND: } On(b,x) \ \land \ Clear(b) \ \land \ Block(b) \ \land \ (b \neq x), \\ \text{EFFECT: } On(b,Table) \ \land \ Clear(x) \ \land \ \neg On(b,x)) \\ \end{array}$

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One solution: [MoveToTable(C, A), Move(B, Table, C), Move(A, Table, B)]

Outline

- Basics on Planning
 - The Problem
 - The PDDL Language
- Search Strategies and Heuristics
 - Forward and Backward Search
 - Heuristics
- Planning Graphs, Heuristics and Graphplan
 - Planning Graphs
 - Heuristics Driven by Planning Graphs
 - The Graphplan Algorithm
- Other Approaches (hints)
 - Planning as SAT Solving
 - Planning as FOL Inference



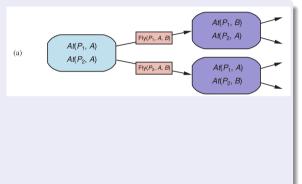
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Two Main Approaches

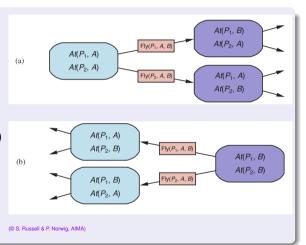
- (a) Forward search (aka progression search)
 - start in the initial state
 - use actions to search forward for a goal state
- (b) Backward search (aka regression search)
 - start from goal states
 - use reverse actions to search forward for the initial state



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- Forward search (aka progression search)
 - choose actions whose preconditions are satisfied
 - add positive effects, delete negative
- Goal test: does the state satisfy the goal?
- Step cost: each action costs 1
- → We can use any of the search algorithms from Ch. 03, 04
 - need keeping track of the actions used to reach the goal
 - Breadth-first and best-first
 - Sound: if they return a plan, then the plan is a solution
 - Complete: if a problem has a solution, then they will return one
 - Require exponential memory wrt. solution length! ⇒ unpractical
 - Depth-first search and greedy search
 - Sound
 - Not complete
 - may enter in infinite loops
 - (classical planning only): made complete by loop-checking
 - Require linear memory wrt. solution length

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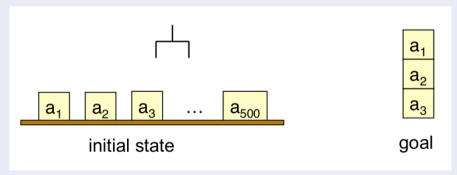
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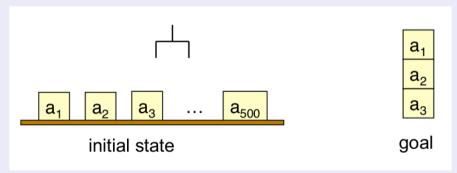
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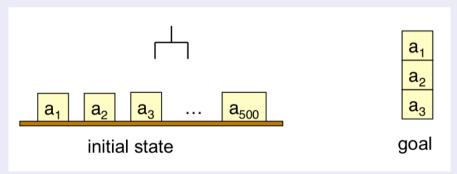
- Planning problems can have huge state spaces
- Forward search can have a very large branching factor
 - ex: pickup(a₁), pickup(a₂), ..., pickup(a₅₀₀)
- \implies Forward-search can waste time trying lots of irrelevant actions
- Need a good heuristic to guide the search



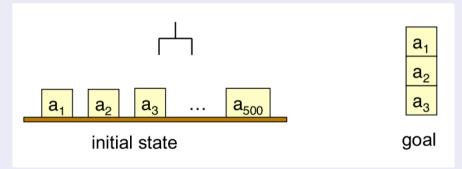
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```
Pos(g') \stackrel{\mathsf{def}}{=} (Pos(g) \setminus Add(a)) \cup Pos(Precond(a))
Neg(g') \stackrel{\mathsf{def}}{=} (Neg(g) \setminus Del(a)) \cup Neg(Precond(a))
```

- Note: Both g and g' represent many states
 - irrelevant ground atoms unassigned
- Consider the goal $At(C_1, SFO) \wedge At(C_2, JFK)$
- Consider the ground action: $Action(Unload(C_1, P_1, SFO), PRECOND : In(C_1, P_1) \land At(P_1, SFO) \land Cargo(C_1) \land Plane(P_1) \land Airport(SFO)$ $EFFECT : At(C_1, SFO) \land \neg In(C_1, P_1))$
- This produces the sub-goal g': $In(C_1, P_1) \wedge At(P_1, SFO) \wedge Cargo(C_1) \wedge Plane(P_1) \wedge Airport(SFO) \wedge At(C_2, JFK)$
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• Predecessor state g' of ground goal g via ground action a:

```
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```
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```

 $PRECOND : In(C_1, P_1) \land At(P_1, SFO) \land Cargo(C_1) \land Plane(P_1) \land Airport(SFO)$ $EFFECT : At(C_1, SFO) \land \neg In(C_1, P_1))$

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```

- Note: Both *q* and *q'* represent many states
 - irrelevant ground atoms unassigned
- Consider the goal $At(C_1, SFO) \wedge At(C_2, JFK)$
- Consider the ground action: Action(Unload(C_1, P_1, SFO).

```
PRECOND: In(C_1, P_1) \land At(P_1, SFO) \land Cargo(C_1) \land Plane(P_1) \land Airport(SFO)
```

 $EFFECT: At(C_1, SFO) \land \neg In(C_1, P_1)$

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```
PRECOND : In(C_1, P_1) \land At(P_1, SFO) \land Cargo(C_1) \land Plane(P_1) \land Airport(SFO)

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- This produces the sub-goal g':
 - $In(C_1, P_1) \wedge At(P_1, SFO) \wedge Cargo(C_1) \wedge Plane(P_1) \wedge Airport(SFO) \wedge At(C_2, JFK)$
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 - e.g. truth value of $In(C_3, P_2)$ irrelevant

- Idea: deal with partially un-instantiated actions and states
 - avoid unnecessary instantiations
 - ⇒ no need to produce a goal for every possible instantiation
- use the most general unifier
- standardize action schemata first (rename vars into fresh ones)
- Consider the goal $At(C_1, SFO) \wedge At(C_2, JFK)$
- Consider the partially-instantiated action: $Action(Unload(C_1, p', SFO), PRECOND: In(C_1, p') \land At(p', SFO) \land Cargo(C_1) \land Plane(p') \land Airport(SFO) \land FFECT: At(C_1, SFO) \land \neg In(C_1, p'))$
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```
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```

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- Represents states with all possible planes
 - \implies no need to produce a subgoal for every plane $P_1, P_2, P_3, ...$

- Idea: deal with partially un-instantiated actions and states
 - avoid unnecessary instantiations
 - → no need to produce a goal for every possible instantiation
- use the most general unifier
- standardize action schemata first (rename vars into fresh ones)
- Consider the goal $At(C_1, SFO) \wedge At(C_2, JFK)$
- Consider the partially-instantiated action:

```
Action(Unload(C_1, p', SFO),
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- Recall: A* is a best-first algorithm which
 - uses an evaluation function f(s) = g(s) + h(s),
 - g(s): (exact) cost to reach s
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Ignore-Preconditions Heuristics

- Ignore all preconditions drops all preconditions from actions
 - every action is applicable in any state
 - any single goal literal can be satisfied in one step (or there is no solution)
 - fast, but over-optimistic
- Remove all preconditions & effects, except literals in the goal
 - more accurate
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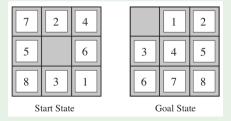
Ignore-Preconditions Heuristics: Example

Sliding tiles

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Action(Slide(t, s_1, s_2),
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 $PRECOND: Tile(t) \land Blank(s_2) \land On(t, s_1) \land Adjacent(s_1, s_2)$ $EFFECT: On(t, s_2) \land Blank(s_1) \land \neg On(t, s_1) \land \neg Blank(s_2)$

- Remove the preconditions $Blank(s_2) \wedge Adjacent(s_1, s_2)$
 - ⇒ we get the number-of-misplaced-tiles heuristics
- Remove the precondition $Blank(s_2)$
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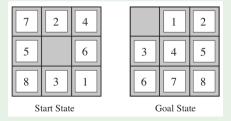
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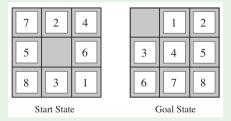
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- Assumption: goals & preconditions contain only positive literals
 - reasonable in many domains
- Idea: Remove the delete lists from all actions
 - No action will ever undo the effect of actions,
 - there is a monotonic progress towards the goal
- Still NP-hard to find the optimal solution of the relaxed problem
 - can be approximated in polynomial time, with hill-climbing
- Can be very effective for some problems

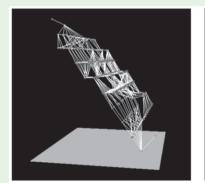
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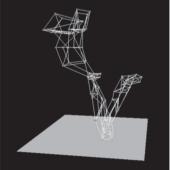
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Ignore Delete-list Heuristics: Example (Hoffmann'05)

- Planning state spaces with ignore-delete-lists heuristic
 - height above the bottom plane is the heuristic score of a state
 - states on the bottom plane are goals
- > No local minima, non dead-ends, non backtracking
- \implies Search for the goal is straightforward for hill-climbing

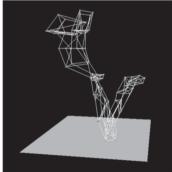




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- Many-to-one mapping from states in the ground/original representation of the problem to a more abstract representation
 - drastically reduces the number of states
- Common strategy: ignore some (less-relevant) fluents
 - drop k fluents \Longrightarrow reduce search space by 2^k factors
 - relevance based on (heuristic) evaluation or domain knowledge
- Air cargo problem: 10 airports, 50 planes, 200 pieces of cargo $\implies 10^{50} \cdot (50 + 10)^{200} \approx 10^{405}$ states (*)
- Consider particular problem in that domain
 - all packages are at 5 airports
 - all packages at a given airport have the same destination
- Abstraction: drop all "At" fluents except for these involving one plane and one package a each of the the 5 airports
 - $\implies 10^5 \cdot (5 + 10)^5 \approx 10^{11} \text{ states (*)}$
 - ullet abstract solution shorter than ground solutions \Longrightarrow admissible
 - abstract solution easy to extend: add Load and Unload actions

(*) wrong in AIMA III Ed, corrected in later editions

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Other Strategies for Planning

Other strategies to define heuristics

- Problem decomposition
 - "divide & conquer" problem into subproblem
 - solve subproblems independently
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 - can drive an algorithm called Graphplan
- A polynomial-size approximation to the (exponential) search tree
 - can be constructed very quickly
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 - can be used to give better heuristic estimates h(s)
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 - level A_i : contains all ground actions with preconditions in S_{i-1}
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 - each S_i may contain both P_j and $\neg P_j$

until
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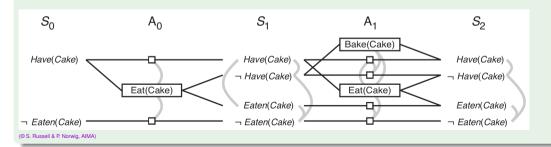
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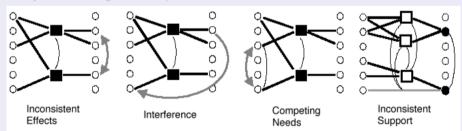
Planning Graph: Example

```
\begin{array}{lll} Init(Have(Cake)) & \textit{You would like to eat your cake and still have a cake.} \\ Goal(Have(Cake) \land Eaten(Cake)) & \textit{Fortunately, you can bake a new one.} \\ Action(Eat(Cake) & \text{Fortunately, you can bake a new one.} \\ & \text{PRECOND: } Have(Cake) & \text{Eaten}(Cake)) \\ & \text{Action}(Bake(Cake) & \text{Small squares persistence actions (no-ops)} \\ & \text{Straight lines indicate preconditions} \\ & \text{and effects} \\ & \text{Effect: } Have(Cake)) & \text{Mutex links are shown as curved gray lines} \\ \end{array}
```



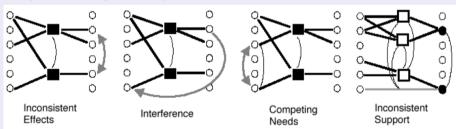
Mutex Computation

- Two actions at the same action-level have a mutex relation if
 - Inconsistent effects: an effect of one negates an effect of the other
 - Interference: one deletes a precondition of the other
 - Inconsistent preconditions (aka competing needs): they have mutually exclusive preconditions
- Otherwise they don't interfere with each other
 - ⇒ both may appear in a solution plan
- Two literals at the same state-level have a mutex relation if
 - inconsistent support: one is the negation of the other
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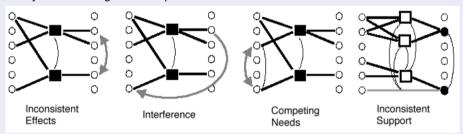
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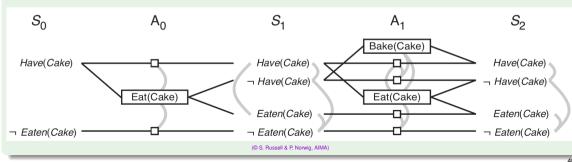
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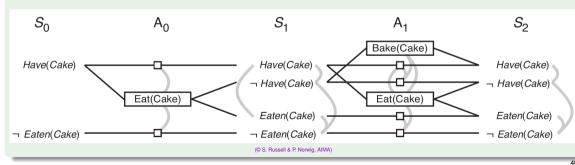


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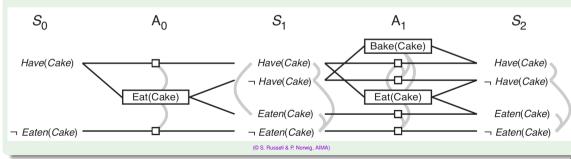
- Inconsistent effects: an effect of one negates an effect of the other ex: persistence of Have(Cake), Eat(Cake) have competing effects ex: Bake(Cake), Eat(Cake) have competing effects
- Interference: one deletes a precondition of the other
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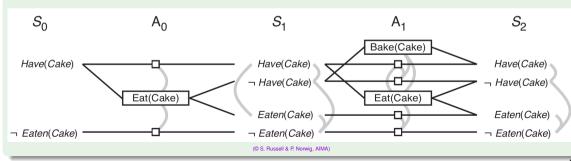
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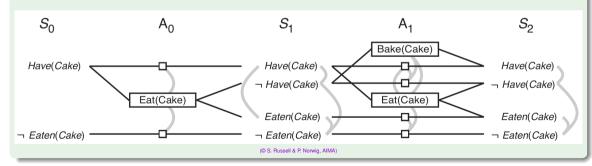


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Mutex Computation: Example [cont.]

- Two literals at the same state-level have a mutex relation if
 - inconsistent support: one is the negation of the other ex.: Have(Cake), ¬Have(Cake)
 - all ways of achieving them are pairwise mutex
 ex.: (S₁): Have(Cake) in mutex with Eaten(Cake)
 because persistence of Have(Cake), Eat(Cake) are mutex



Create initial layer S_0 :

 \bigcirc insert into S_0 all literals in the initial state

Repeat for increasing values of i = 0, 1, 2, ...

Create action layer Ai

- of for each action schema, for each way to unify its preconditions to non-mutually exclusive literals in
- \bigcirc for every literal in S_i , enter a no-op action node into A_i
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Planning Graphs: Complexity

- A planning graph is polynomial in the size of the problem:
 - a graph with n levels, a actions, I literals, has size $O(n(a+l)^2)$
 - time complexity is also $O(n(a+l)^2)$
- ⇒ The process of constructing the planning graph is very fast
 - does not require choosing among actions

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 - The Problem
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- Each level S_i represents a set of possible belief states
 - two literals connected by a mutex belong to different belief state
- A literal not appearing in the final level of the graph cannot be achieved by any plan
 - \implies if a goal literal is not in the final level, the problem is unsolvable
- The level S_j a literal I appears first is never greater than the level it can be achieved in a plan
 - *j* is called the level cost of literal /
- the level cost of a literal g_i in the graph constructed starting from state s, is an estimate of the cost to achieve it from s (i.e. h(g))
 - this estimate is admissible
 - ex: from s₀ Have(cake) has cost 0 and Eaten(cake) has cost 1
- Planning graph admits several actions per level
 - ⇒ inaccurate estimate
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Estimating the heuristic cost of a conjunction of goal literals

- Max-level heuristic: the maximum level cost of the sub-goals
 - admissible
- Level-sum heuristic: the sum of the level costs of the goals
 - inadmissible only if goals are independent,
 - it may work well in practice
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Planning Graphs for Heuristic Estimation [cont.]

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```

- A strategy for extracting a plan from the planning graph
- Repeatedly adds a level to a planning graph (EXPAND-GRAPH)
- If all the goal literals occur in last level and are non-mutex
 - search for a plan that solves the problem (EXTRACT-SOLUTION)
 - if that fails, expand another level and try again (and add $\langle goal, level \rangle$ as nogood)
- If graph and nogoods have both leveled off then return failure
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[Recall] Example: Spare Tire Problem

```
Init(Tire(Flat) \land Tire(Spare) \land At(Flat, Axle) \land At(Spare, Trunk))
Goal(At(Spare, Axle))
Action(Remove(obj, loc),
  PRECOND: At(obj, loc)
  EFFECT: \neg At(obj, loc) \land At(obj, Ground)
Action(PutOn(t, Axle),
   PRECOND: Tire(t) \land At(t, Ground) \land \neg At(Flat, Axle)
   EFFECT: \neg At(t, Ground) \land At(t, Axle)
Action(LeaveOvernight,
   PRECOND:
   EFFECT: \neg At(Spare, Ground) \land \neg At(Spare, Axle) \land \neg At(Spare, Trunk)
            \wedge \neg At(Flat, Ground) \wedge \neg At(Flat, Axle) \wedge \neg At(Flat, Trunk))
```

(© S. Russell & P. Norwig, AIMA)

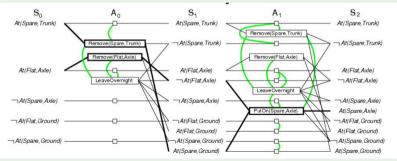
(We assume that the car is parked in a particularly bad neighborhood, so that the effect of leaving it overnight is that the tires disappear.)

One solution: [Remove(Flat, Axle), Remove(Spare, Trunk), PutOn(Spare, Axle)]

Graphplan: Example

Spare Tire problem

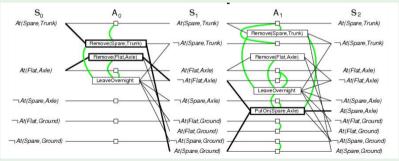
- Initial plan 5 literals from initial state and the Closed-World-Assumption literals (S_0).
 - fixed literals (e.g. Tire(Flat)) ignored here
 - irrelevant literals ignored here
- Goal At(Spare, Axle) not present in S₀
 - ⇒ no need to call EXTRACT-SOLUTION
- ullet Graph and nogoods not leveled off \Longrightarrow invoke EXPAND-GRAPH



Graphplan: Example [cont.]

Spare Tire problem

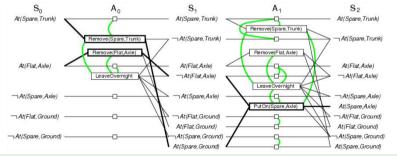
- Invoke EXPAND-GRAPH
 - add actions A₀, persistence actions and mutexes
 - add fluents S₁ and mutexes
- Goal At(Spare, Axle) not present in S₁
 - ⇒ no need to call EXTRACT-SOLUTION
- ullet Graph and nogoods not leveled off \Longrightarrow invoke EXPAND-GRAPH



Graphplan: Example [cont.]

Spare Tire problem

- Invoke Expand-Graph
 - add actions A₁, persistence actions and mutexes
- add fluents S₂ and mutexes
- Goal At(Spare, Axle) present in S₂
 - call Extract-Solution
- Solution found!



Exercise

- Consider the following variant of the Spare Tire problem: add At(Flat, Trunk) to the goal
- Write the (non-serialized) planning graph
- Extract a plan from the graph
- Do the same with the serialized planning graph

Graphplan "family" of algorithms, depending on approach used in EXTRACT-SOLUTION(...)

- Can be formulated as an (incremental) SAT problem
 - one proposition for each ground action and fluent
 - clauses represent preconditions, effects, no-ops and mutexes
- Can be formulated as a backward search problem
- Planning problem restricted to planning graph
 - mutexes found by EXPAND-GRAPH prune paths in the search tree
 - → much faster than unrestricted planning
- (if P.G. not serialized) may produce partial order plans
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Partial-Order Plans

Partial-Order vs. Total-Order Plans

- Total-order plans: strictly linear sequences of actions
 - disregards the fact that some action are mutually independent
- Partial-order plans: set of precedence constraints between action pairs
 - form a directed acyclic graph
 - longest path to goal may be much shorter than total-order plan
 - easily converted into (possibly many) distinct total-order plans (any possible interleaving of independent actions)

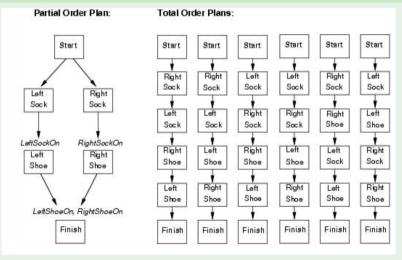
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Partial-Order Plans: Example

Socks & Shoes Examples



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Termination of Graphplan

- Theorem: If the graph and the no-goods have both leveled off, and no solution is found we can safely terminate with failure
- Intuition (proof sketch):
 - Literals and actions increase monotonically and are finite
 - Mutex and no-goods decrease monotonically (and cannot become less than zero)
 - ⇒ they too eventually must level off
 - When we reach this stable state, if one of the goals is missing or is mutex with another goal, then it will remain so
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Exercise

- Socks & Shoes example:
 - Formalize the Socks & Shoes example in PDDL
 - Write the non-serialized planning graph
 - Compute the level cost for every fluent
 - Ohoose some states, compute h(s) using the three heuristics
 - Extract a plan from the graph in (2)
 - Ompare h(s) with the level they occur in the plan
 - Write the serialized planning graph
 - Repeat steps (3)-(6) with the serialized graph
- Do same steps (1)-(8) for the Air Cargo Transport example

Outline

- Basics on Planning
 - The Problem
 - The PDDL Language
- Search Strategies and Heuristics
 - Forward and Backward Search
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- Encode bounded planning problem into a propositional formula
- ⇒ Solve it by (incremental) calls to a SAT solver
 - A model for the formula (if any) is a plan of length t
 - Many variants in the encoding
 - Extremely efficient with many problems of interest

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Planning as SAT Solving [cont.]

- TRANSLATE-TO-SAT(INIT, TRANSITION, GOAL, T):
 - ground fluents & actions at each step are propositionalized
 - ex: $\langle At(P_1, SFO), 3 \rangle \Longrightarrow At_P_1_SFO_3$
 - ex: $\langle Fly(P_1, SFO, JFK), 3 \rangle \Longrightarrow Fly_P_1_SFO_JFK_3$
 - returns propositional formula: $Init^0 \wedge (\bigwedge_{i=1}^{t-1} Transition^{i,i+1}) \wedge Goal^t$
- Init⁰ and Goal^t: conjunctions of literals at step 0 and t resp.
 - ex: $Init^0$: $At_P_1_SFO_0 \land At_P_2_JFK_0$
 - ex: $Goal^3$: At_{P_1} JFK_3 \wedge At_{P_2} SFO_3
- *Transition*ⁱ, i+1</sup>: encodes transition from steps i to i+1
 - Actions: $Action^i \rightarrow (Precond^i \land Effects^{i+1})$ ex: $Fly_P_1_SFO_JFK_2 \rightarrow (At_P_1_SFO_2 \land At_P_1_JFK_3$
 - No-Ops: for each fluent *F* and step *i*:

$$F^{i+1} \leftrightarrow \bigvee ActionCausingF_k^i \lor (F^i \land \bigwedge \lnot ActionCausingNotF_j^i)$$

- Mutex constraints: ¬Action¹₁ ∨ ¬Action¹₂
 ex: ¬Fly_P₁_SFO_JFK_2 ∨ ¬Fly_P₁_SFO_Newark_2
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 - ex: $Init^0$: At P_1 SFO $0 \land At$ P_2 JFK 0
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Exercise

Consider the socks & shoes example

- Trenslate it into SAT for t=0,1,2
 - non serialized
 - no need to propositionalize: treat ground atoms as propositions
 - no need to CNF-ize here (human beings don't like CNFs)
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- Idea: formalize planning into FOL
- ⇒ use resolution-based inference for planning
 - + Admit quantifications --> very expressive
 - allows formalizing sentences like "move all the cargos from A to B regardless of how many pieces of cargo there are"
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- Situation:
 - the initial state is a situation
 - if s is a situation and a is an action, then Result(s, a) is a situation
 - Result() injective: Result(s, a) = Result(s', a') \leftrightarrow ($s = s' \land a = a'$)
 - a solution is a situation that satisfies the goal
- ullet Action preconditions: $\Phi(s) o Possible(a, s)$
 - \bullet $\Phi(s)$ describes preconditions
 - $\bullet \;\; \mathsf{ex} \colon (\mathit{Alive}(\mathit{Agent}, s) \land \mathit{Have}(\mathit{Agent}, \mathit{Arrow}, s)) \rightarrow \mathit{Possible}(\mathit{Shoot}, s)$
- Successor-state axioms (similar to propositional case):

$$[Action is possible] \rightarrow \begin{bmatrix} [Fluent is true in result state] \leftrightarrow \\ ([Action's effect made it true] \lor \\ ([It was true before] \land [action left it alone])) \end{bmatrix}$$

- ex: $Possible(a, s) \rightarrow \left[\begin{array}{c} Holding(Agent, g, Result(a, s)) \leftrightarrow \\ a = Grab(g) \lor (Holding(Agent, g, s) \land a \ne Release(g)) \end{array} \right]$
- Unique action axioms: $A_i(x,...) \neq A_i(y,...)$ ex $Shoot(x) \neq Grab(y)$
- A_i injective: $A_i(x_1,...,x_n) = A_i(y_1,...,y_n) \leftrightarrow \bigwedge_{i=1}^n x_i = y_i$, ex: $Grab(x) = Grab(y) \leftrightarrow x = y$

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 - if s is a situation and a is an action, then Result(s, a) is a situation
 - Result() injective: Result(s, a) = Result(s', a') \leftrightarrow (s = s' \land a = a')
 - a solution is a situation that satisfies the goal
- Action preconditions: $\Phi(s) \rightarrow Possible(a, s)$
 - Φ(s) describes preconditions
 - ex: $(Alive(Agent, s) \land Have(Agent, Arrow, s)) \rightarrow Possible(Shoot, s)$
- Successor-state axioms (similar to propositional case):

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[Action is possible] \rightarrow \begin{bmatrix} [Fluent is true in result state] \leftrightarrow \\ ([Action's effect made it true] \lor \\ ([It was true before] \land [action left it alone])) \end{bmatrix}
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- ex: $Possible(a, s) \rightarrow \left[\begin{array}{c} Holding(Agent, g, Result(a, s)) \leftrightarrow \\ a = Grab(g) \lor (Holding(Agent, g, s) \land a \ne Release(g)) \end{array} \right]$
- Unique action axioms: $A_i(x,...) \neq A_i(y,...)$ ex $Shoot(x) \neq Grab(y)$
- A_i injective: $A_i(x_1,...,x_n) = A_i(y_1,...,y_n) \leftrightarrow \bigwedge_{i=1}^n x_i = y_i$, ex: $Grab(x) = Grab(y) \leftrightarrow x = y$

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Situation Calculus: Example

Situations as the results of actions in the Wumpus world

