

Mathematical Logic

First order logic: syntax and semantics

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November 10, 2014

- Why First Order Logic (FOL)?
- Syntax and Semantics of FOL;
- First Order Theories;
- ... and in between few examples;

Question

Try to express in Propositional Logic the following statements:

- Mary is a person
- John is a person
- Mary is mortal
- Mary and John are siblings

A solution

Through atomic propositions:

- Mary-is-a-person
- John-is-a-person
- Mary-is-mortal
- Mary-and-John-are-siblings

Problem with previous solution

- Mary-is-a-person
- John-is-a-person
- Mary-is-mortal
- Mary-and-John-are-siblings

How do we link Mary of the first sentence to Mary of the third sentence? Same with John. How do we link Mary and Mary-and-John?

Question

Try to express in Propositional Logic the following statements:

- All persons are mortal;
- There is a person who is a spy.

A solution

We can give all people a name and express this fact through atomic propositions:

- $\text{Mary-is-mortal} \wedge \text{John-is-mortal} \wedge \text{Chris-is-mortal} \wedge \dots \wedge \text{Michael-is-mortal}$
- $\text{Mary-is-a-spy} \vee \text{John-is-a-spy} \vee \text{Chris-is-a-spy} \vee \dots \vee \text{Michael-is-a-spy}$

Problem with previous solution

- $\text{Mary-is-mortal} \wedge \text{John-is-mortal} \wedge \text{Chris-is-mortal} \wedge \dots \wedge \text{Michael-is-mortal}$
- $\text{Mary-is-a-spy} \vee \text{John-is-a-spy} \vee \text{Chris-is-a-spy} \vee \dots \vee \text{Michael-is-a-spy}$

The representation is not compact and generalization patterns are difficult to express.

What if we do not know all the people in our “universe”? How can we express the statement independently from the people in the “universe”?

Question

Try to express in Propositional Logic the following statements:

- Every natural number is either even or odd

A solution

We can use two families of propositions $even_i$ and odd_i for every $i \geq 1$, and use the set of formulas

$$\{odd_i \vee even_i \mid i \geq 1\}$$

$$\{odd_i \vee even_i \mid i \geq 1\}$$

What happens if we want to state this in one single formula? To do this we would need to write an infinite formula like:

$$(odd_1 \vee even_1) \wedge (odd_2 \vee even_2) \wedge \dots$$

and this cannot be done in propositional logic.

Question

Express the statements:

- the father of Luca is Italian

Solution (Partial)

- `mario-is-father-of-luca` \supset `mario-is-italian`
- `michele-is-father-of-luca` \supset `michele-is-italian`
- ...

Problem with previous solution

- `mario-is-father-of-luca` \supset `mario-is-italian`
- `michele-is-father-of-luca` \supset `michele-is-italian`
- ...

This statement strictly depend from a fixed set of people. What happens if we want to make this statement independently of the set of persons we have in our universe?

Why first order logic?

Because it provides a way of **representing** information like the following one:

- 1 **Mary is a person;**
- 2 John is a person;
- 3 **Mary is mortal;**
- 4 Mary and John are siblings
- 5 **Every person is mortal;**
- 6 There is a person who is a spy;
- 7 Every natural number is either even or odd;
- 8 The father of Luca is Italian

and also to **infer** the third one from the first one and the fifth one.

Whereas propositional logic assumes world contains facts, first-order logic (like natural language) assumes the world contains:

- **Constants:** mary, john, 1, 2, 3, red, blue, world war 1, world war 2, 18th Century...
- **Predicates:** Mortal, Round, Prime, Brother of, Bigger than, Inside, Part of, Has color, Occurred after, Owns, Comes between, ...
- **Functions:** Father of, Best friend, Third inning of, One more than, End of, ...

Constants and Predicates

- Mary is a person
- John is a person
- Mary is mortal
- Mary and John are siblings

In FOL it is possible to build atomic propositions by applying a **predicate** to **constants**

- *Person(mary)*
- *Person(john)*
- *Mortal(mary)*
- *Siblings(mary, john)*

Quantifiers and variables

- Every person is mortal;
- There is a person who is a spy;
- Every natural number is either even or odd;

In FOL it is possible to build propositions by applying **universal** (**existential**) **quantifiers** to **variables**. This allows to quantify to arbitrary objects of the universe.

- $\forall x. Person(x) \supset Mortal(x)$;
- $\exists x. Person(x) \supset Spy(x)$;
- $\forall x. (Odd(x) \vee Even(x))$

- The father of Luca is Italian.

In FOL it is possible to build propositions by applying a **function** to a **constant**, and then a predicate to the resulting object.

- *Italian(fatherOf(Luca))*

Syntax of FOL

The **alphabet of FOL** is composed of two sets of symbols:

Logical symbols

- the logical constant \perp
- propositional logical connectives $\wedge, \vee, \supset, \neg, \equiv$
- the **quantifiers** \forall, \exists
- an infinite set of **variable symbols** x_1, x_2, \dots
- the **equality symbol** $=$. (optional)

Non Logical symbols

- a set c_1, c_2, \dots of **constant symbols**
- a set f_1, f_2, \dots of **functional symbols** each of which is associated with its **arity** (i.e., number of arguments)
- a set P_1, P_2, \dots of **relational symbols** each of which is associated with its **arity** (i.e., number of arguments)

Non logical symbols - Example

Non logical symbols depends from the domain we want to model.
Their must have an intuitive interpretation on such a domain.

Example (Domain of arithmetics)

symbols	type	arity	intuitive interpretation
0	constant	0*	the smallest natural number
$\text{succ}(\cdot)$	function	1	the function that given a number returns its successor
$+(\cdot, \cdot)$	function	2	the function that given two numbers returns the number corresponding to the sum of the two
$<(\cdot, \cdot)$	relation	2	the less then relation between natural numbers

* A constant can be considered as a function with arity equal to 0

Non logical symbols - Example

Example (Domain of arithmetics - extended)

The basic language of arithmetics can be extended with further symbols e.g:

symbols	type	arity	intuitive interpretation
0	constant	0	the smallest natural number
$succ(\cdot)$	function	1	the function that given a number returns its successor
$+(\cdot, \cdot)$	function	2	the function that given two numbers returns the number corresponding to the sum of the two
$*(\cdot, \cdot)$	function	2	the function that given two numbers returns the number corresponding to the product of the two
$<(\cdot, \cdot)$	relation	2	the less then relation between natural numbers
$\leq(\cdot, \cdot)$	relation	2	the less then or equal relation between natural numbers

Non logical symbols - Example

Example (Domain of strings)

symbols	type	arity	intuitive interpretation
ϵ	constant	0	The empty string
"a", "b",	constants	0	The strings containing one single character of the latin alphabet
$concat(\cdot, \cdot)$	function	2	the function that given two strings returns the string which is the concatenation of the two
$subst(\cdot, \cdot, \cdot)$	function	3	The function that replaces all the occurrence of a string with another string in a third one
$<$	relation	2	Alphabetic order on the strings
$substring(\cdot, \cdot)$	relation	2	a relation that states if a string is contained in another string

Terms and formulas of FOL

Terms

- every constant c_i and every variable x_i is a term;
- if t_1, \dots, t_n are terms and f_i is a functional symbol of arity equal to n , then $f(t_1, \dots, t_n)$ is a term

Well formed formulas

- if t_1 and t_2 are terms then $t_1 = t_2$ is a formula
- If t_1, \dots, t_n are terms and P_i is relational symbol of arity equal to n , then $P_i(t_1, \dots, t_n)$ is formula
- if A and B are formulas then $\perp, A \wedge B, A \supset B, A \vee B, \neg A, A \equiv B$ are formulas
- if A is a formula and x a variable, then $\forall x.A$ and $\exists x.A$ are formulas.

Example (Terms)

- x_i ,
- c_i ,
- $f_i(x_j, c_k)$, and
- $f(g(x, y), h(x, y, z), y)$

Example (formulas)

- $f(a, b) = c$,
- $P(c_1)$,
- $\exists x(A(x) \vee B(y))$, and
- $P(x) \supset \exists y.Q(x, y)$.

An example of representation in FOL

Example (Language)

constants	functions (arity)	Predicate (arity)
Aldo	mark (2)	attend (2)
Bruno	best-friend (1)	friend (2)
Carlo		student (1)
MathLogic		course (1)
DataBase		less-than (2)
0, 1, ..., 10		

Example (Terms)

Intuitive meaning

an individual named Aldo
the mark 1
Bruno's best friend
anything
Bruno's mark in MathLogic
somebody's mark in DataBase
Bruno's best friend mark in MathLogic

term

Aldo
1
best-friend(Bruno)
x
mark(Bruno,MathLogic)
mark(x,DataBase)
mark(best-friend(Bruno),MathLogic)

An example of representation in FOL (cont'd)

Example (Formulas)

Intuitive meaning	Formula
Aldo and Bruno are the same person	$Aldo = Bruno$
Carlo is a person and MathLogic is a course	$person(Carlo) \wedge course(MathLogic)$
Aldo attends MathLogic	$attend(Aldo, MathLogic)$
Courses are attended only by students	$\forall x(attend(x, y) \supset course(y) \supset student(x))$
every course is attended by somebody	$\forall x(course(x) \supset \exists y attend(y, x))$
every student attends something	$\forall x(student(x) \supset \exists y attend(x, y))$
a student who attends all the courses	$\exists x(student(x) \wedge \forall y(course(y) \supset attend(x, y)))$
every course has at least two attenders	$\forall x(course(x) \supset \exists y \exists z(attend(y, x) \wedge attend(z, x) \wedge \neg y = z))$
Aldo's best friend attend the same courses attended by Aldo	$\forall x(attend(Aldo, x) \supset attend(best\text{-}friend(Aldo), x))$
best-friend is symmetric	$\forall x(best\text{-}friend(best\text{-}friend(x)) = x)$
Aldo and his best friend have the same mark in MathLogic	$mark(best\text{-}friend(Aldo), MathLogic) = mark(Aldo, MathLogic)$
A student can attend at most two courses	$\forall x \forall y \forall z \forall w(attend(x, y) \wedge attend(x, z) \wedge attend(x, w) \supset (y = z \vee z = w \vee y = w))$

Common Mistakes

- Use of \wedge with \forall

$\forall x (WorksAt(FBK, x) \wedge Smart(x))$ means “Everyone works at FBK and everyone is smart”

“Everyone working at FBK is smart” is formalized as
 $\forall x (WorksAt(FBK, x) \supset Smart(x))$

- Use of \supset with \exists

$\exists x (WorksAt(FBK, x) \supset Smart(x))$ means “There is a person so that if (s)he works at FBK then (s)he is smart” and this is true as soon as there is at last an x who does not work at FBK

“There is an FBK-working smart person” is formalized as
 $\exists x (WorksAt(FBK, x) \wedge Smart(x))$

Representing variations quantifiers in FOL

Example

Represent the statement **at least 2** students attend the KR course

$$\exists x_1 \exists x_2 (\text{attend}(x_1, KR) \wedge \text{attend}(x_2, KR))$$

The above representation is not enough, as x_1 and x_2 are variable and they could denote the same individual, we have to guarantee the fact that x_1 and x_2 denote different person. The correct formalization is:

$$\exists x_1 \exists x_2 (\text{attend}(x_1, KR) \wedge \text{attend}(x_2, KR) \wedge x_1 \neq x_2)$$

At least n ...

$$\exists x_1 \dots x_n \left(\bigwedge_{i=1}^n \phi(x_i) \wedge \bigwedge_{i \neq j=1}^n x_i \neq x_j \right)$$

Example

Represent the statement **at most 2** students attend the KR course

$$\forall x_1 \forall x_2 \forall x_3 (attend(x_1, KR) \wedge attend(x_2, KR) \wedge attend(x_3, KR) \supset x_1 = x_2 \vee x_2 = x_3 \vee x_1 = x_3)$$

At most n ...

$$\forall x_1 \dots x_{n+1} \left(\bigwedge_{i=1}^{n+1} \phi(x_i) \supset \bigvee_{i \neq j=1}^{n+1} x_i = x_j \right)$$

FOL interpretation for a language L

A first order interpretation for the language

$L = \langle c_1, c_2, \dots, f_1, f_2, \dots, P_1, P_2, \dots \rangle$ is a pair $\langle \Delta, \mathcal{I} \rangle$ where

- Δ is a non empty set called **interpretation domain**
- \mathcal{I} is a function, called **interpretation function**
 - $\mathcal{I}(c_i) \in \Delta$ (elements of the domain)
 - $\mathcal{I}(f_i) : \Delta^n \rightarrow \Delta$ (n -ary function on the domain)
 - $\mathcal{I}(P_i) \subseteq \Delta^n$ (n -ary relation on the domain)

where n is the arity of f_i and P_i .

Example of interpretation

Example (Of interpretation)

Symbols Constants: *alice*, *bob*, *carol*, *robert*
Function: *mother-of* (with arity equal to 1)
Predicate: *friends* (with arity equal to 2)

Domain $\Delta = \{1, 2, 3, 4, \dots\}$

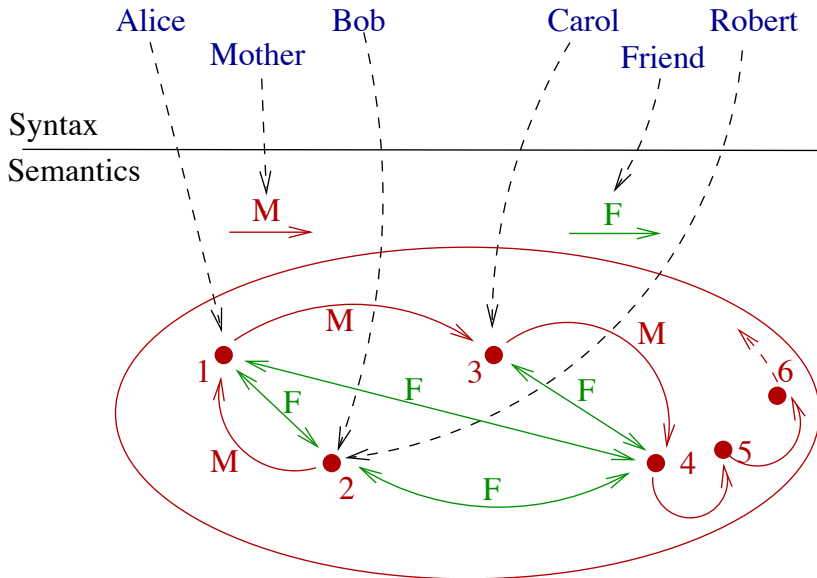
Interpretation $\mathcal{I}(\textit{alice}) = 1$, $\mathcal{I}(\textit{bob}) = 2$, $\mathcal{I}(\textit{carol}) = 3$,
 $\mathcal{I}(\textit{robert}) = 2$

$\mathcal{I}(\textit{mother-of}) = M$

$M(1) = 3$
 $M(2) = 1$
 $M(3) = 4$
 $M(n) = n + 1$ for $n \geq 4$

$\mathcal{I}(\textit{friends}) = F = \left\{ \begin{array}{l} \langle 1, 2 \rangle, \langle 2, 1 \rangle, \langle 3, 4 \rangle, \\ \langle 4, 3 \rangle, \langle 4, 2 \rangle, \langle 2, 4 \rangle, \\ \langle 4, 1 \rangle, \langle 1, 4 \rangle, \langle 4, 4 \rangle \end{array} \right\}$

Example (cont'd)



Interpretation of terms

Definition (Assignment)

An **assignment** a is a function from the set of variables to Δ .

$a[x/d]$ denotes the assignment that coincides with a on all the variables but x , which is associated to d .

Definition (Interpretation of terms)

The **interpretation** of a term t w.r.t. the assignment a , in symbols $\mathcal{I}(t)[a]$ is recursively defined as follows:

$$\mathcal{I}(x_i)[a] = a(x_i)$$

$$\mathcal{I}(c_i)[a] = \mathcal{I}(c_i)$$

$$\mathcal{I}(f(t_1, \dots, t_n))[a] = \mathcal{I}(f)(\mathcal{I}(t_1)[a], \dots, \mathcal{I}(t_n)[a])$$

FOL Satisfiability of formulas

Definition (Satisfiability of a formula w.r.t. an assignment)

An interpretation \mathcal{I} **satisfies** a formula ϕ w.r.t. the assignment a according to the following rules:

$$\mathcal{I} \models t_1 = t_2[a] \quad \text{iff} \quad \mathcal{I}(t_1)[a] = \mathcal{I}(t_2)[a]$$

$$\mathcal{I} \models P(t_1, \dots, t_n)[a] \quad \text{iff} \quad \langle \mathcal{I}(t_1)[a], \dots, \mathcal{I}(t_n)[a] \rangle \in \mathcal{I}(P)$$

$$\mathcal{I} \models \phi \wedge \psi[a] \quad \text{iff} \quad \mathcal{I} \models \phi[a] \text{ and } \mathcal{I} \models \psi[a]$$

$$\mathcal{I} \models \phi \vee \psi[a] \quad \text{iff} \quad \mathcal{I} \models \phi[a] \text{ or } \mathcal{I} \models \psi[a]$$

$$\mathcal{I} \models \phi \supset \psi[a] \quad \text{iff} \quad \mathcal{I} \not\models \phi[a] \text{ or } \mathcal{I} \models \psi[a]$$

$$\mathcal{I} \models \neg\phi[a] \quad \text{iff} \quad \mathcal{I} \not\models \phi[a]$$

$$\mathcal{I} \models \phi \equiv \psi[a] \quad \text{iff} \quad \mathcal{I} \models \phi[a] \text{ iff } \mathcal{I} \models \psi[a]$$

$$\mathcal{I} \models \exists x\phi[a] \quad \text{iff} \quad \text{there is a } d \in \Delta \text{ such that } \mathcal{I} \models \phi[a[x/d]]$$

$$\mathcal{I} \models \forall x\phi[a] \quad \text{iff} \quad \text{for all } d \in \Delta, \mathcal{I} \models \phi[a[x/d]]$$

Exercise

Check the following statements, considering the interpretation \mathcal{I} defined few slides ago:

- 1 $\mathcal{I} \models \textit{Alice} = \textit{Bob}[a]$
- 2 $\mathcal{I} \models \textit{Robert} = \textit{Bob}[a]$
- 3 $\mathcal{I} \models x = \textit{Bob}[a[x/2]]$

Example (cont'd)

$$\mathcal{I}(\text{mother-of}(\text{alice}))[a] = 3$$

$$\mathcal{I}(\text{mother-of}(x))[a[x/4]] = 5$$

$$\mathcal{I}(\text{friends}(x, y)) =$$

$x :=$	$y :=$
1	2
2	1
4	1
1	4
4	2
2	4
4	3
3	4
4	4

$$\mathcal{I}(\text{friends}(x, x)) =$$

$x :=$
4

$$\mathcal{I}(\text{friends}(x, y) \wedge x = y) =$$

$x :=$	$y :=$
4	4

$$\mathcal{I}(\exists x \text{friends}(x, y)) =$$

$y :=$
2
1
4
3

$$\mathcal{I}(\forall x \text{friends}(x, y)) =$$

$y :=$
4

Analogy with Databases

When the language \mathcal{L} and the domain of interpretation Δ are finite, and \mathcal{L} does not contains functional symbols (relational language), there is a strict **analogy between first order logics and databases**.

- Non logical simbols of \mathcal{L} correspond to database schema (tables)
- Δ corresponds to the set of values which appears in the tables (active domain)
- the interpretation \mathcal{I} corresponds to the tuples that belongs to each relation
- Formulas on \mathcal{L} corresponds to query over the database
- Interpretation of formulas of \mathcal{L} correspond to answers.

Analogy with Databases

FOL	DB
<i>friends</i>	CREATE TABLE FRIENDS (friend1 : INTEGER friend2 : INTEGER)
<i>friends</i> (x, y)	SELECT * FROM FRIENDS
<i>friends</i> (x, x)	SELECT friend1 FROM FRIENDS WHERE friends1 = friends2
<i>friends</i> (x, y) $\wedge x = y$	SELECT * FROM FRIENDS WHERE friends1 = friends2
$\exists x.friends(x, y)$	SELECT friend2 FROM FRIENDS
<i>friends</i> (x, y) \wedge <i>friends</i> (y, z)	SELECT * FROM FRIENDS as FRIEND1 FRIENDS as FRIEND2 WHERE FRIENDS1.friends2 = FRIENDS2.friends1

Free variables

Intuition

A **free occurrence** of a variable x is an occurrence of x which is not bounded by a (universal or existential) quantifier.

Definition (Free occurrence)

- any occurrence of x in t_k is free in $P(t_1, \dots, t_k, \dots, t_n)$
- any free occurrence of x in ϕ or in ψ is also free in $\phi \wedge \psi$, $\psi \vee \phi$, $\psi \supset \phi$, and $\neg\phi$
- any free occurrence of x in ϕ , is free in $\forall y.\phi$ and $\exists y.\phi$ if y is distinct from x .

Definition (Ground/Closed Formula)

A formula ϕ is **ground** if it does not contain any variable. A formula is **closed** if it does not contain free occurrences of variables.

A **variable x is free** in ϕ (denote by $\phi(x)$) if there is at least a free occurrence of x in ϕ .

Free variables represents individuals which must be instantiated to make the formula a meaningful proposition.

- x is free in $friends(alice, x)$.
- x is free in $P(x) \supset \forall x.Q(x)$ (the occurrence of x in red is free the one in green is not free).

Intuitively..

Free variables represents individuals which must be instantiated to make the formula a meaningful proposition.

- $Friends(Bob, y)$ y free
- $\forall y.Friends(Bob, y)$ no free variables
- $Sum(x, 3) = 12$ x free
- $\exists x.(Sum(x, 3) = 12)$ no free variables
- $\exists x.(Sum(x, y) = 12)$ y free

Definition (Term free for a variable)

A term t is free for a variable x in formula ϕ , if and only if all the occurrences of x in ϕ do not occur within the scope of a quantifier of some variable occurring in t .

Example

The term x is free for y in $\exists z.hates(y, z)$. We can safely replace y with x obtaining $\exists z.hates(x, z)$ without changing the meaning of the formula.

However, the term z is not free for y in $\exists z.hates(y, z)$. In fact y occurs within the scope of a quantifier of z . Thus, we cannot substitute z for y in this sentence without changing the meaning of the sentence as we obtain $\exists z.hates(z, z)$.

Free variables and free terms - example

An occurrence of a variable x can be safely instantiated by a **term free for x in a formula ϕ** ,

If you replace x with a terms which is not free for x in ϕ , you can have unexpected effects:

E.g., replacing x with *mother-of*(y) in the formula $\exists y.\textit{friends}(x, y)$ you obtain the formula

$$\exists y.\textit{friends}(\textit{mother-of}(y), y)$$

Satisfiability and Validity

Definition (Model, satisfiability and validity)

An interpretation \mathcal{I} is a **model** of ϕ under the assignment a , if

$$\mathcal{I} \models \phi[a]$$

A formula ϕ is **satisfiable** if there is some \mathcal{I} and some assignment a such that $\mathcal{I} \models \phi[a]$.

A formula ϕ is **unsatisfiable** if it is not satisfiable.

A formula ϕ is **valid** if every \mathcal{I} and every assignment a $\mathcal{I} \models \phi[a]$

Definition (Logical Consequence)

A formula ϕ is a **logical consequence** of a set of formulas Γ , in symbols $\Gamma \models \phi$, if for all interpretations \mathcal{I} and for all assignment a

$$\mathcal{I} \models \Gamma[a] \quad \implies \quad \mathcal{I} \models \phi[a]$$

where $\mathcal{I} \models \Gamma[a]$ means that \mathcal{I} satisfies all the formulas in Γ under a .

Say where these formulas are valid, satisfiable, or unsatisfiable

- $\forall x P(x)$
- $\forall x P(x) \supset \exists y P(y)$
- $\forall x. \forall y. (P(x) \supset P(y))$
- $P(x) \supset \exists y P(y)$
- $P(x) \vee \neg P(y)$
- $P(x) \wedge \neg P(y)$
- $P(x) \supset \forall x. P(x)$
- $\forall x \exists y. Q(x, y) \supset \exists y \forall x Q(x, y)$
- $x = x$
- $\forall x. P(x) \equiv \forall y. P(y)$
- $x = y \supset \forall x. P(x) \equiv \forall y. P(y)$
- $x = y \supset (P(x) \equiv P(y))$
- $P(x) \equiv P(y) \supset x = y$

$\forall xP(x)$	Satisfiable
$\forall xP(x) \supset \exists yP(y)$	Valid
$\forall x.\forall y.(P(x) \supset P(y))$	Satisfiable
$P(x) \supset \exists yP(y)$	Valid
$P(x) \vee \neg P(y)$	Satisfiable
$P(x) \wedge \neg P(y)$	Satisfiable
$P(x) \supset \forall x.P(x)$	Satisfiable
$\forall x\exists y.Q(x, y) \supset \exists y\forall xQ(x, y)$	Satisfiable
$x = x$	Valid
$\forall x.P(x) \equiv \forall y.P(y)$	Valid
$x = y \supset \forall x.P(x) \equiv \forall y.P(y)$	Valid
$x = y \supset (P(x) \equiv P(y))$	Valid
$P(x) \equiv P(y) \supset x = y$	Satisfiable

Properties of quantifiers

Proposition

The following formulas are valid

- $\forall x(\phi(x) \wedge \psi(x)) \equiv \forall x\phi(x) \wedge \forall x\psi(x)$
- $\exists x(\phi(x) \vee \psi(x)) \equiv \exists x\phi(x) \vee \exists x\psi(x)$
- $\forall x\phi(x) \equiv \neg\exists x\neg\phi(x)$
- $\forall x\exists x\phi(x) \equiv \exists x\phi(x)$
- $\exists x\forall x\phi(x) \equiv \forall x\phi(x)$

Proposition

The following formulas are not valid

- $\forall x(\phi(x) \vee \psi(x)) \equiv \forall x\phi(x) \vee \forall x\psi(x)$
- $\exists x(\phi(x) \wedge \psi(x)) \equiv \exists x\phi(x) \wedge \exists x\psi(x)$
- $\forall x\phi(x) \equiv \exists x\phi(x)$
- $\forall x\exists y\phi(x, y) \equiv \exists y\forall x\phi(x, y)$

Expressing properties in FOL

What is the meaning of the following FOL formulas?

- 1 $\exists x(bought(Frank, x) \wedge dvd(x))$
 - 2 $\exists x.bought(Frank, x)$
 - 3 $\forall x.(bought(Frank, x) \rightarrow bought(Susan, x))$
 - 4 $(\forall x.bought(Frank, x)) \rightarrow (\forall x.bought(Susan, x))$
 - 5 $\forall x \exists y.bought(x, y)$
 - 6 $\exists x \forall y.bought(x, y)$
-
- 1 "Frank bought a dvd."
 - 2 "Frank bought something."
 - 3 "Susan bought everything that Frank bought."
 - 4 "If Frank bought everything, so did Susan."
 - 5 "Everyone bought something."
 - 6 "Someone bought everything."

Expressing properties in FOL

Define an appropriate language and formalize the following sentences using FOL formulas.

- 1 All Students are smart.
- 2 There exists a student.
- 3 There exists a smart student.
- 4 Every student loves some student.
- 5 Every student loves some other student.
- 6 There is a student who is loved by every other student.
- 7 Bill is a student.
- 8 Bill takes either Analysis or Geometry (but not both).
- 9 Bill takes Analysis and Geometry.
- 10 Bill doesn't take Analysis.
- 11 No students love Bill.

Expressing properties in FOL

- 1 $\forall x.(Student(x) \rightarrow Smart(x))$
- 2 $\exists x.Student(x)$
- 3 $\exists x.(Student(x) \wedge Smart(x))$
- 4 $\forall x.(Student(x) \rightarrow \exists y.(Student(y) \wedge Loves(x, y)))$
- 5 $\forall x.(Student(x) \rightarrow \exists y.(Student(y) \wedge \neg(x = y) \wedge Loves(x, y)))$
- 6 $\exists x.(Student(x) \wedge \forall y.(Student(y) \wedge \neg(x = y) \rightarrow Loves(y, x)))$
- 7 $Student(Bill)$
- 8 $Takes(Bill, Analysis) \leftrightarrow \neg Takes(Bill, Geometry)$
- 9 $Takes(Bill, Analysis) \wedge Takes(Bill, Geometry)$
- 10 $\neg Takes(Bill, Analysis)$
- 11 $\neg \exists x.(Student(x) \wedge Loves(x, Bill))$

Expressing properties in FOL

For each property write a formula expressing the property, and for each formula write the property it formalises.

- Every Man is Mortal

$$\forall x. Man(x) \supset Mortal(x)$$

- Every Dog has a Tail

$$\forall x. Dog(x) \supset \exists y (PartOf(x, y) \wedge Tail(y))$$

- There are two dogs

$$\exists x, y (Dog(x) \wedge Dog(y) \wedge x \neq y)$$

- Not every dog is white

$$\neg \forall x. Dog(x) \supset White(x)$$

- $\exists x. Dog(x) \wedge \exists y. Dog(y)$

There is a dog

- $\forall x, y (Dog(x) \wedge Dog(y) \supset x = y)$

There is at most one dog

Open and Closed Formulas

- Note that for closed formulas, satisfiability, validity and logical consequence do not depend on the assignment of variables.
- For closed formulas, we therefore omit the assignment and write $\mathcal{I} \models \phi$.
- More in general $\mathcal{I} \models \phi[a]$ if and only if $\mathcal{I} \models \phi[a']$ when $[a]$ and $[a']$ coincide on the variables free in ϕ (they can differ on all the others)

(un)satisfiability/validity of a FOL formula - examples

Example

Decide whether or not $\forall x(P(x) \supset Q(x)) \supset (\forall xP(x) \supset \forall xQ(x))$ is valid.

- The above formula is valid when $\mathcal{I} \models \forall x(P(x) \supset Q(x)) \supset (\forall xP(x) \supset \forall xQ(x))[a]$ for all assignment a . Which is equivalent to say that
- if $\mathcal{I} \models \forall x(P(x) \supset Q(x))[a]$ then $\mathcal{I} \models (\forall xP(x) \supset \forall xQ(x))[a]$; which is the same as:
- if $\mathcal{I} \models \forall x(P(x) \supset Q(x))[a]$ and $\mathcal{I} \models \forall xP(x)[a]$ then $\mathcal{I} \models \forall xQ(x)[a]$.
- To show the previous fact, suppose that:
(H1) $\mathcal{I} \models \forall x(P(x) \supset Q(x))[a]$, and that
(H2) $\mathcal{I} \models \forall xP(x)[a]$.
- From the hypothesis (H1), we have that for all $d \in \Delta^{\mathcal{I}}$,
 $\mathcal{I} \models P(x) \supset Q(x)[a[x/d]]$
- from the hypothesis (H2), we have that for all $d \in \Delta^{\mathcal{I}}$, $\mathcal{I} \models P(x)[a[x/d]]$
- by the definition of satisfiability of implication we have that for all $d \in \Delta^{\mathcal{I}}$,
 $\mathcal{I} \models Q(x)[a[x/d]]$
- which implies that $\mathcal{I} \models \forall xQ(x)[a]$.

(un)satisfiability/validity of a FOL formula - examples

Example

Check if the formula $(\forall xP(x) \supset \forall xQ(x)) \supset \forall x(P(x) \supset Q(x))$ is valid:

- This time we try to show that the formula is **not valid**.
- For this we have to find an interpretation \mathcal{I} such that $\mathcal{I} \models \forall xP(x) \supset \forall xQ(x)[a]$ but $\mathcal{I} \not\models \forall x(P(x) \supset Q(x))[a]$.
- in order to have that $\mathcal{I} \models \forall xP(x) \supset \forall xQ(x)[a]$, we can choose to falsify the premise of the implication, i.e., to build an interpretation such that $\mathcal{I} \not\models \forall xP(x)[a]$.
- we need an element d in the domain of interpretation $\Delta^{\mathcal{I}}$, such that $\mathcal{I} \not\models P(x)[a[x/d]]$.
- In order to have that $\mathcal{I} \not\models \forall x(P(x) \supset Q(x))[a]$, we need an element d' of the domain $\Delta^{\mathcal{I}}$ such that $\mathcal{I} \models P(x)[a[x/d']]$ and $\mathcal{I} \not\models Q(x)[a[x/d']]$.
- at this point we can build the interpretation \mathcal{I} on the domain $\Delta^{\mathcal{I}} = \{d, d'\}$ with $P^{\mathcal{I}} = \{d'\}$ and $Q^{\mathcal{I}} = \emptyset$.

Exercise (FOL Theory)

Let \mathcal{L} be a first order language on a signature containing

- the constant symbols a and b ,
- the binary function symbol f , and
- the binary predicate symbol P .

Answer to the following questions:

- 1 Does \mathcal{L} have a finite model? If yes define it, if not explain why.
- 2 Let \mathcal{T} be a theory containing the following axioms
 - 1 $\forall y. \neg P(x, x)$ (P is irreflexive)
 - 2 $\forall xyz. (P(x, y) \wedge P(y, z) \supset P(x, z))$ (P is transitive)
 - 3 $\forall xy. (P(x, f(x, y)) \wedge P(y, f(x, y)))$

Is \mathcal{T} satisfiable?. If yes can you provide a model for \mathcal{T}

- 3 Does \mathcal{T} have a finite model? If yes, define it; if not, explain why.

Exercise (FOL Theory)

Suppose that a first order language L contains only the set of constants $\{a, b, c\}$ and no functional symbols, and the unary predicate symbol P .

Say if the following formula is valid, i.e., true in all interpretations. If it is valid give a proof of its validity; if it is not valid provide a counter-model.

$$P(a) \wedge P(b) \wedge P(c) \supset \forall x P(x)$$

Exercise (FOL Modelling)

Transform in FOL the following sentences:

- 1 ciao ciao
- 2 The fathers of dogs are dogs.
- 3 There are at least two students enrolled in every course.
- 4 No region is part of each of two disjoint regions

Transform in Natural Language the following sentences:

- 1 $\forall x(Bag(x) \supset \exists y(Coin(y) \wedge Contains(x, y)))$
- 2 $\exists x(Telephone(x) \wedge \forall y(Secretary(y) \supset \neg Uses(x, y)))$
- 3 $\exists x(Buyer(x) \wedge Bought(x, TheScream) \wedge \forall y(Buyer(y) \wedge Bought(y, TheScream) \supset x = y))$