

# Computational Linguistics: Semantics II

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# 1. Recall: Formal Semantics Main questions

The main questions are:

1. What does a given sentence mean?
2. How is its meaning built?
3. How do we infer some piece of information out of another?

## 1.1. Building Meaning Representations

To build a meaning representation we need to fulfill three tasks:

**Task 1** Specify a reasonable **syntax** for the natural language fragment of interest.

**Task 2** Specify semantic representations for the **lexical items**.

**Task 3** Specify the **translation** of constituents **compositionally**. That is, we need to specify the translation of such expressions in terms of the translation of their parts, parts here referring to the substructure given to us by the syntax.

Moreover, when interested in Computational Semantics, all three tasks need to be carried out in a way that leads to computational implementation naturally.

## 1.2. Lambda-calculus: Functional Application

Summing up:

- ▶ FA has the form:  $\text{Functor}(\text{Argument})$ . E.g.  $(\lambda x.\text{love}(x, \text{mary}))(\text{john})$
- ▶ FA triggers a very simple operation: Replace the  $\lambda$ -bound variable by the argument. E.g.  $(\lambda x.\text{love}(x, \text{mary}))(\text{john}) \Rightarrow \text{love}(\text{john}, \text{mary})$

▶ “Every student left”

$$\forall x.\text{Student}(x) \rightarrow \text{left}(x)$$

▶ “A student left”

$$\exists x.\text{Student}(x) \wedge \text{left}(x)$$

▶ “No student left”

$$\neg \exists x.\text{Student}(x) \wedge \text{left}(x)$$

But how do we reach these meaning representations starting from the lexicon?

## 2.2. Determiners (cont'd)

Let's start representing "a man" as  $\exists x.man(x)$ . Applying the rules we have seen so far, we obtain that the representation of "A man loves Mary" is:

$$love(\exists x.man(x), mary)$$

which is clearly wrong.

Notice that  $\exists z.man(z)$  just isn't the meaning of "a man". If anything, it translates the complete sentence "There is a man".



## 2.3. Determiners (Cont'd)

Let's start from what we have, namely “man’ and “loves Mary”:

$\lambda y.man(y)$ ,  $\lambda x.love(x, mary)$ .

Hence, the term representing “a” is:

$$\lambda X.\lambda Y.\exists z.X(z) \wedge Y(z)$$

Try to obtain the meaning representation for “a man”, and the “a man loves Mary”.

By  $\beta$ -conversion twice we obtain that “a man” is  $\lambda Y.\exists z.Man(z) \wedge Y(z)$ , and then  $\exists z.Man(z) \wedge love(z, mary)$

### 3. Ambiguities

Starting from:

john: $j$	book: $\lambda x(\mathbf{book}(x))$
read: $\lambda x.\lambda y.(\mathbf{read}(y, x))$	didn't: $\lambda X.\neg X$
a: $\lambda X.\lambda Y(\exists x.X(x) \wedge Y(x))$	

Built the meaning representation for “John didn’t read a book”.

a.  $\exists x.\mathbf{book}(x) \wedge \neg\mathbf{read}(j, x)$  [A > NOT]

b.  $\neg\exists x.B(x) \wedge \mathbf{read}(j, x)$  [NOT > A]

► **Scope:** In a. the quantifier phrase (QP), “a book”, has scope over “didn’t” [A > NOT], whereas in b. it has narrow scope [NOT > A].

► **Binding:** the variable  $x$  is bound by “a book” in “John didn’t read a book”.

## 3.1. Scope Ambiguities

Can you think of other expressions that may cause scope ambiguity?

John **think** a student left

Does the student exist or not?

a.  $\exists x.think(j, left(x))$

b.  $think(j, \exists x.left(x))$

## 4. Dependencies

While studying the syntax of natural language, we have seen that important concepts to account for are local and long-distance dependencies.

The  $\lambda$ -operator gives us (more or less) a way to represent this link semantically.

For instance, in  $\lambda x.\lambda y.like(y, x)$  we express that the dependency of the subject and object from the verb.

But the calculus gives us also a natural way to handle long-distance dependencies: eg. relative pronouns.

## 4.1. Relative Pronouns

For instance, “which John read [...]”:

We know how to represent the noun phrase “John” and the verb “read”, namely, as `john` and `λx.y.read(y, x)`.

What is the role of “which” in e.g. “the book which John read is read”?

The term representing “which” has to express the fact that it is replacing the role of a noun phrase in subject (or object position) within a subordinate sentence while being the subject (object) of the main sentence:

$$\lambda X.\lambda Y.\lambda z.X(z) \wedge Y(z)$$

The double role of “which” is expressed by the double occurrence of  $z$ .

## 4.2. Relative Pronoun (Cont'd)

Recall,

$$\lambda X.\lambda Y.\lambda z.X(z) \wedge Y(z)$$

- i. read u:  $\lambda y(\text{read}(y, u))$       ii. John read u:  $\text{read}(j, u)$   
iii. John read:  $\lambda u.\text{read}(j, u)$       iv. which John read:  $\lambda Y.\lambda z.\text{read}(j, z) \wedge Y(z)$

- ▶ at the syntactic level we said that the relative pronoun “which” plays the role of the verb’s object and it leaves a **gap** in the object position.
- ▶ Semantically, the gap is represented by the  $u$  on which the relative pronoun forces the abstraction [iii.] before taking its place.

## 5. Summing up: Constituents and Assembly

Let's go back to the points where FOL fails, i.e. constituent representation and assembly. The  $\lambda$ -calculus succeeds in both:

**Constituents:** each constituent is represented by a lambda term.

John:  $j$  knows:  $\lambda xy.(\text{know}(x))(y)$  read john:  $\lambda y.\text{know}(y, j)$

**Assembly:** function application ( $\alpha(\beta)$ ) and abstraction ( $\lambda x.\alpha[x]$ ) capture **composition** and **decomposition** of meaning representations.

## 6. Lambda-Terms Interpretations

In the Logic course you've seen that an Interpretation is a pair consisting of a domain ( $\mathcal{D}$ ) and an interpretation function ( $\mathcal{I}$ ).

- ▶ In the case of FOL we had only one domain, namely the one of the objects/entities we were reasoning about. Similarly, we only had one type of variables. Moreover, we were only able to speak of propositions/clauses.
- ▶  $\lambda$ -terms speak of functions and we've used also **variables standing for functions**. Therefore, we need a more complex concept of interpretation, or better a more **complex concept of domain** to provide the fine-grained distinction among the objects we are interested in: truth values, entities and functions.
- ▶ For this reason, the  $\lambda$ -calculus is of Higher Order.



## 6.1. Models, Domains, Interpretation

In order to interpret meaning representations expressed in FOL augmented with  $\lambda$ , the following facts are essential:

- ▶ **Sentences:** Sentences can be thought of as referring to their truth value, hence they denote in the the domain  $D_t = \{1, 0\}$ .
- ▶ **Entities:** Entities can be represented as constants denoting in the domain  $D_e$ , e.g.  $D_e = \{\text{john}, \text{vincent}, \text{mary}\}$
- ▶ **Functions:** The other natural language expressions can be seen as incomplete sentences and can be interpreted as **boolean functions** (i.e. functions yielding a truth value). They denote on functional domains  $D_b^{D_a}$  and are represented by functional terms of type  $(a \rightarrow b)$ .

For instance “walks” misses the subject (of type  $e$ ) to yield a sentence ( $t$ ).

- ▶ denotes in  $D_t^{D_e}$
- ▶ is of type  $(e \rightarrow t)$ ,
- ▶ is represented by the term  $\lambda x_e(\text{walk}(x))_t$

## 6.2. Lambda-calculus: some remarks

The pure lambda calculus is a theory of functions as rules invented around 1930 by Church. It has more recently been applied in Computer Science for instance in “Semantics of Programming Languages”.

In Formal Linguistics we are mostly interested in lambda conversion and abstraction. Moreover, we work only with typed-lambda calculus and even more, only with a fragment of it.

The types are the ones we have seen above labeling the domains, namely:

- ▶  $e$  and  $t$  are types.
- ▶ If  $a$  and  $b$  are types, then  $(a \rightarrow b)$  is a type.

## 7. The Three Tasks Revised

**Task 1** Specify a reasonable **syntax** for the natural language fragment of interest.

**We can do this using CFG.**

**Task 2** Specify semantic representations for the **lexical items**. **We know what this involves**

**Task 3** Specify the **translation** of an item  $\mathcal{R}$  whose parts are  $\mathcal{F}$  and  $\mathcal{A}$  with the help of functional application. That is, we need to specify which part is to be thought of as functor (here it's  $\mathcal{F}$ ), which as argument (here it's  $\mathcal{A}$ ) and then let the resultant translation  $\mathcal{R}'$  be  $\mathcal{F}'(\mathcal{A}')$ . **We know that  $\beta$ -conversion (with the help of  $\alpha$ -conversion), gives us the tools needed to actually construct the representation built by this process.**

In the Lab we have seen that the solution proposed for task 1 leads itself to computational implementation naturally. Next week we will see that this holds for task 2 and 3 too (though we won't go into the detail of it. If you are interested in it and you know Prolog already: good topic for a project!)

## 8. Historical Remark: Montague

Montague ('74) was the first to seriously propose and defend the thesis that the relation between syntax and semantics in a natural language such as English could be viewed as not essentially different from the relation between syntax and semantics in formal language such as the language of FOL. The framework he developed is known as “Montague’s Universal Grammar” and its main components are:

- ▶ Model-Theory
- ▶ The principle of “Compositionality” which is due to Frege (1879).
- ▶ The  $\lambda$ -calculus which was invented by Church in the 1930’s.
- ▶ Categorical Grammar which is due to Ajdukiewicz ('35) and Bar-Hill ('53). It’s a (context free) Grammar based on functional application.

The novelty of Montagues’ work was to apply them to natural language in a uniform framework.

We will study Categorical Grammar later in the course.

## 9. Inference

We've said that an important ultimate question in Formal Semantics is “How do we infer some piece of information out of other?”

We have seen how to use logic to represent natural language input.

Therefore the question reduces to the question of “Which are the inference tools for the logic we used?”.

For instance, a tool you already know is the “Tableaux Method” (**Logic** Course by Franconi).

In the **Computational Logic** course (by Filottrani) you will study the computational properties of this tool and get hands on experience implementing it in PROLOG.

## 10. Inference (Cont'd)

Inference play a crucial role also at the discourse level, which is one of the next topics. For instance,

“Jon has a rabbit. The tail is white and fluffy”.

“The tail” of whom?

I know that most rabbits have a tail, hence Jon’s rabbit has a tail. Therefore, I can interpret “the tail” to be of Jon’s rabbit.

## 11. The need for semantic representation

We shall also assume a framework in which we are attempting to communicate in natural language with some computer system which has its own internal representation of “knowledge” in some symbolic form.

The natural language front end (NLFE) has to formulate queries, commands and statements in a way that allow appropriate responses to be produced by the “back-end” system.

This can be done in two ways, using:

- ▶ Syntactic Representation
- ▶ Meaning Representation

## 11.1. Syntactic Representation

In this framework, it would be possible in principle to have the syntactic form of an input sentence (i.e. the syntax tree) act directly upon the stored information within the back-end knowledge based system, without any intermediate representations being constructed. That is, the response of the system, whether it be the moving of a robot arm, or the printing out of information, would be computed directly from the syntax tree.

This approach is **rarely followed**, as it leads to unnecessary complications. There are certain regularities about how each natural language expresses information, and there may be several different ways of expressing essentially the same request. If the syntactic tree is directly interpreted, then all the (linguistic) knowledge about these **grammatical variations** (which are irrelevant to the work of the back-end) has to be built-in to the interpreter which controls the back-end responses. Minor alterations to the grammatical coverage of the front-end then necessitate changes to this response-generator.



## 11.2. Meaning Representation

It is therefore normal to have the NLFE formulate its query (or command, or whatever) in some specialised notation for semantic representation, and then pass this on to further levels of the system for processing.

The superficial details of how the request was expressed in English can then be handled by separate linguistic rules, without the back-end having to include such information – the semantic representation of the input is a “canonical form”. (This is a very similar methodological argument to that which supports splitting the earlier phase of processing into two stages, syntax and semantic interpretation).

Much of the research that goes on in NL semantics is concerned with the design of suitable formalisms for this intermediate language.

## 11.3. Meaning Representation (Cont'd)

It is important to note that this intermediate semantic language need not be the same as whatever representation the actual back-end uses to store its knowledge. The semantic representation language is suitable for depicting the content of a sentence (or a sequence of sentences), and for extracting a response from the back-end; the back-end's own formalism, used for storing information, might be completely different. This point is sometimes glossed over in systems which use some very general notation (e.g. logic, Prolog) for both purposes.

## 11.4. Remark

Notice that a system of semantic representation is not just some diagrams or symbolic expressions thrown together in a pleasing or intuitively attractive way. Very simple examples such as representing “John loves Mary” by  $love(john, mary)$  can sometimes tempt beginners into thinking that a semantic expression is just the words rearranged, with a bit of punctuation thrown in. This is not a viable approach. A semantic representation formalism should have at the very least the following properties:

1. It should be **clearly defined**. That is, there should be an explicit statement of what would qualify as a legal (well-formed) expression in this notation.
2. Its operations/behaviour should be clearly defined. That is, there should be proper definitions of what **inference** or other manipulations can be carried out on these expressions.
3. To justify the term “semantic”, it should be reasonable to argue that the behaviour of these expressions in some way captures the meaning of the corresponding linguistic expressions.

Anything that lacks these fundamental properties cannot be a serious semantic representation, although it might be handy as an informal notation for sketching out one's thoughts, or a visual aid in conveying ideas to other workers.

## 12. Semantics vs. Knowledge Representation

To a large extent, the issues involved in semantics are the same as the more general issues of knowledge representation.

But the two endeavours (semantics and knowledge representation) are not wholly identical, since knowledge representation could well be less constrained in what is acceptable as a solution.

- ▶ For natural language **semantics**, the formalisations chosen should make those **distinctions which are reflected in the natural language involved**, whereas general knowledge representation need make only those distinctions relevant to inference and other processes.
- ▶ The representations used in general **knowledge representation need not be expressible in natural language at all**, but those in semantics have expressibility in language as their entire justification.
- ▶ More particularly, NL semantic representations (and their rules) must have a well-defined interface to actual words, phrases and sentences (the concrete syntax), which is not a requirement for more general knowledge representation.

## 13. Conclusions