1 PROMELA overview

- Processes
- Data objects
- Message Channels
- Executability



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- Channels are used to transfer messages between active processes.
- They store messages in first-in first-out order.
- Two types:
 - buffered channels,
 - rendezvous ports, also called synchronous channels.

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• Declaration:

```
chan qname = [16] of { short, byte, bool }
```

This channel can store up to 16 messages, each consisting of 3 fields of the types listed.

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qname!expr1,expr2,expr3

The process blocks if the channel is full.

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• Receiving a message:

qname?var1,var2,var3

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The process blocks if the channel is empty.

 Useful pre-defined functions len, empty, nempty, full, nfull: len(qname)

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• The first message field is a message type indication: qname!expr1(expr2,expr3) qname?var1(var2,var3)

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- The first message field is a message type indication: qname!expr1(expr2,expr3) qname?var1(var2,var3)
- Some parameters can be given as constants:

qname?cons1,var2,cons2

The process blocks if the channel is empty or if the sent values do not match the constants.

Rendezvous Ports

Declaration of a rendezvous port (it pass single byte messages)
 chan port = [0] of { byte }

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Rendezvous Ports

- Declaration of a rendezvous port (it pass single byte messages)
 chan port = [0] of { byte }
- The channel size is zero: the channel port can pass, but can not store messages!

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Rendezvous Ports

- Declaration of a rendezvous port (it pass single byte messages)
 chan port = [0] of { byte }
- The channel size is zero: the channel port can pass, but can not store messages!
- Message interaction is synchronous: two processes execute a send and a receive statement at the same time (as a single atomic operation).

```
mtype = { msgtype };
chan name = [0] of { mtype, byte };
active proctype A()
{            name!msgtype(124);
            name!msgtype(121)
}
active proctype B()
{            byte state;
            name?msgtype(state)
```

}

-

Channels of channels

- Message parameters are always passed by value.
- We can pass the value of a channel from a process to another.

```
mtype = { msgtype };
chan glob = [0] of { chan };
active proctype A()
{
        chan loc = [0] of { mtype, byte };
        glob!loc;
        loc?msgtype(121)
}
active proctype B()
{
        chan who;
        glob?who;
        who!msgtype(121)
```

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Statements

- Every statement is either *executable* or *blocked*.
- Three main types of statements:
 - print statements
 - assignments
 - expression statements
- Print statements and assignments are always executable (as well as skip, assert, ...).
- Expression statements are executable iff they evaluate to true.
 - (2 < 3) always executable;
 - (x < 27) blocked until x is less than 27;
 - (3 + x) executable when x differs from -3.
- Expressions must be side effect free

(e.g. b = c++ is not valid).

- Exception: the run statement can be considered as a blocking expression:
 - it blocks when there are 255 processes alive;
 - if it does not block, it creates a new process.

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Given the following PROMELA program:

```
active proctype P() {
    int x = 0;
    x++;
    int y = x;
    assert(y == 1);
}
```

Is the assertion invalid? If yes, why?

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Is the assertion invalid? If yes, why?

All variable declarations are always implicitly moved to the beginning of process.

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Basic verification 2

```
chan com = [0] of { byte };
byte value;
bool d;
proctype p() {
 byte i;
 do
  :: if
     :: i >= 5 -> break
     :: else -> printf("Doing something else\n"); i ++
     fi
  :: com ? value; printf("p received: %d\n",value)
  od;
 d = 1
}
init {
     atomic {
        run p();
     }
     end: com ! 100;
}
```

Is it possible that process p does not read from the channel at all?

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Is it possible that process p does not read from the channel at all? Yes

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Write a PROMELA model for summing up an array of integers.

- Declare and (nondeterministically) initialize an integer array.
- Add a loop that sums up the elements.

Declare a rendezvous channel and create two processes:

- The first process sends the numbers 0 through 9 onto the channel.
- The second process reads the values of the channel and outputs them.

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Declare a rendezvous channel and create two processes:

- The first process sends the numbers 0 through 9 onto the channel.
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Replace the rendezvous with a buffered channel and check how the behavior changes.