Spin: Overview of PROMELA*

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Contents

- PROMELA overview
 - Processes
 - Data objects
 - Message Channels

2 Exercises

PROMELA

PROMELA is not a programming language, but rather a meta-language for building verification models.

- The design of PROMELA is focused on the interaction among processes at the system level;
- Provides:
 - non-deterministic control structures,
 - primitives for process creation,
 - primitives for interprocess communication.
- Misses:
 - functions with return values,
 - expressions with side-effects,
 - data and functions pointers.

Types of objects

Three basic types of objects:

- processes
- data objects
- message channels

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Process Initialization [1/3]

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- active: process created at initialization phase (right after init)
 active [2] proctype you_run() {
 printf("my pid is: %d\n", _pid)
 }

Process Initialization [1/3]

- init is always the first process initialized (if declared);
- active: process created at initialization phase (right after init)

```
active [2] proctype you_run() {
         printf("my pid is: %d\n", _pid)
}
```

• run: process created when instruction is processed

note: run allows for input parameters!

Process Initialization [2/3]

No parameter can be given to init nor to active processes.

```
active proctype proc(byte x) {
   printf("x = %d\n", x);
}
```

```
• "$ spin test.pml x = 0
```

Process Initialization [2/3]

No parameter can be given to init nor to active processes.

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If present, active process parameters default to 0.
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Process Initialization [2/3]

No parameter can be given to init nor to active processes.

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active proctype proc(byte x) {
    printf("x = %d\n", x);
}
If present, active process parameters default to 0.
```

A process does not necessarily start right after being created

```
proctype proc(byte x) {
    printf("x = %d\n", x);
}
init {
    run proc(0);
    run proc(1);
}
* spin test.pml

x = 0

x = 1

x = 0

x = 1

x = 0
```

Process Initialization [3/3]

Only a limited number of processes (255) can be created:

Process Initialization [3/3]

Only a limited number of processes (255) can be created:

- A process "terminates" when it reaches the end of its code.
- A process "dies" when it has terminated and all processes created after it have died.

Process Execution [1/2]

- Processes execute **concurrently** with all other processes.
- Processes are scheduled non-deterministically.
- Processes are interleaved: statements of different processes do not occur at the same time (except for synchronous channels).
- Each process may have several different possible actions enabled at each point of execution: only one choice is made (non-deterministically).

Process Execution [2/2]

- Each process has its own local state:
 - process counter _pid (location within the proctype);
 - value of the local variables.
- A process communicates with other processes:
 - using global (shared) variables (might need synchronization!);
 - using channels.

Statements [1/2]

- each statement is atomic
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 - print statements
 - assignments
 - skip
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 - break
 - ...

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- each statement is atomic
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- Always executable:
 - print statements
 - assignments
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 - break
 - ...
- Not always executable:
 - the run statement is executable only if there are less than 255 processes alive;
 - expressions

Statements [2/2]

- An expression is executable iff it evaluates to true (i.e. non-zero).
 - (5 < 30): always executable;
 - (x < 30): blocks if x is not less than 30;
 - (x + 30): blocks if x is equal to -30;
- Busy-Waiting: the expression (a == b); is equivalent to: while (a != b) { skip }; /* C-code */
- Expressions must be side-effect free (e.g. b = c++ is not valid).

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Basic types

Туре	Typical Range
bit	0, 1
bool	false, true
byte	0255
chan	1255
mtype	1255
pid	0255
short	$-2^{15} \dots 2^{15} - 1$
int	$-2^{31} \cdot \cdot 2^{31} - 1$
unsigned	$02^{n}-1$

- A byte can be printed as a character with the %c format specifier;
- There are no floats and no strings;

Typical declarations

```
/* two single bits, initially 0
bit x, y;
                                                                    */
bool turn = true:
                               /* boolean value, initially true
                                                                    */
byte a[12];
                               /* all elements initialized to 0
                                                                    */
byte a[3] = {'h', 'i', '\setminus 0'};
                               /* byte array emulating a string
                                                                    */
chan m:
                               /* uninitialized message channel
                                                                    */
mtype n;
                               /* uninitialized mtype variable
                                                                    */
short b[4] = 89;
                               /* all elements initialized to 89
                                                                    */
                                                                    */
int cnt = 67;
                               /* integer scalar, initially 67
unsigned v : 5;
                               /* unsigned stored in 5 bits
                                                                    */
unsigned w : 3 = 5;
                               /* value range 0..7, initially 5
                                                                    */
```

- All variables are initialized by default to 0.
- Array indexes starts at 0.

Data structures

 A run statement accepts a list of variables or structures, but no array.

```
typedef Record {
    byte a[3];
    int x;
    bit b
};
proctype run_me(Record r) {
    r.x = 12
}
init {
    Record test;
    run run_me(test)
}
```

Note: but array can still be enclosed in data structures

 Multi-dimensional arrays are not supported, although there are indirect ways:

```
typedef Array {
          byte el[4]
};
Array a[4];
```

Variable Scope

- Spin (old versions): only two levels of scope
 - global scope: declaration outside all process bodies.
 - local scope: declaration within a process body.

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- Spin (old versions): only two levels of scope
 - global scope: declaration outside all process bodies.
 - local scope: declaration within a process body.
- Spin (versions 6+): added block-level scope

Note: since Spin version 2.0, variable declarations are not implicitly moved to the beginning of a block

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Message Channels

- A channel is a FIFO (first-in first-out) message queue.
- A channel can be used to exchange messages among processes.
- Two types:
 - buffered channels,
 - synchronous channels (aka rendezvous ports)

 Declaration of a channel storing up to 16 messages, each consisting of 3 fields of the listed types:

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chan qname = [16] of { short, byte, bool }
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 Note: array must be enclosed within user-defined types.
- Useful pre-defined functions: len, empty, nempty, full, nfull: len(qname);

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qname!expr1,expr2,expr3
```

The process blocks if the channel is full.

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- Message Send:

```
qname!expr1,expr2,expr3
```

The process blocks if the channel is full.

• Message Receive:

```
qname?var1,var2,var3
```

The process blocks if the channel is empty.

Alternative use of Buffered Channels

An alternative syntax for message send/receive involves brackets:

```
qname!expr1(expr2,expr3)
qname?var1(var2,var3)
```

It can be used to highlight that the first message field is interpreted as 'message type'.

Alternative use of Buffered Channels

- An alternative syntax for message send/receive involves brackets:
 qname!expr1(expr2,expr3)
 qname?var1(var2,var3)
 It can be used to highlight that the first message field is interpreted as 'message type'.
- If at the receiving side some parameter is set to a constant value: qname?const1,var2,var3
 then the process blocks if the channel is empty or the input message field does not match the fixed constant value.

Synchronous Channels

• A synchronous channel (aka rendezvous port) has size zero.

```
chan port = [0] of { byte }
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- A synchronous channel (aka rendezvous port) has size zero.chan port = [0] of { byte }
- Messages can be exchanged, but not stored!
- Synchronous execution: a process executes a send at the same time another process executes a receive (as a single atomic operation).

Channels of channels

- Message parameters are always passed by value.
- We can also pass the value of a channel from a process to another.

```
mtype = { msgtype };
chan glob = [0] of { chan };
active proctype A() {
        chan loc = [0] of { mtype, byte };
        glob!loc; /* send channel loc through glob */
        loc?msgtype(121) /* read 121 from channel loc
                                                           */
}
active proctype B() {
        chan who;
        glob?who;
                          /* receive channel loc from glob */
        who!msgtype(121) /* write 121 on channel loc
}
```

Q: what if B sends 122 on channel loc?

Channels of channels

- Message parameters are always passed by value.
- We can also pass the value of a channel from a process to another.

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mtype = { msgtype };
chan glob = [0] of { chan };
active proctype A() {
        chan loc = [0] of { mtype, byte };
        glob!loc; /* send channel loc through glob */
        loc?msgtype(121) /* read 121 from channel loc
                                                           */
}
active proctype B() {
        chan who;
        glob?who;
                          /* receive channel loc from glob */
        who!msgtvpe(121)
                          /* write 121 on channel loc
                                                           */
}
```

Q: what if B sends 122 on channel loc? both A and B are forever blocked

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Basic verification

```
chan com = [0] of { byte };
byte value;
proctype p() {
    byte i;
    dο
        :: if
               :: i >= 5 -> break
               :: else -> printf("Doing something else\n"); i ++
           fi
        :: com ? value; printf("p received: %d\n",value)
    od;
    ... /* fill in for formal verification */
init {
    run p();
    end: com ! 100;
}
```

 \mathbf{Q} : is it possible that process \mathbf{p} does not read from the channel at all?

Basic verification

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        :: com ? value; printf("p received: %d\n",value)
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    ... /* fill in for formal verification */
init {
    run p();
    end: com ! 100;
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```

 $\boldsymbol{Q}\boldsymbol{:}$ is it possible that process \boldsymbol{p} does not read from the channel at all? Yes

Exercises

- Ex. 1: write a PROMELA model that sums up an array of integers.
 - declare and (non-deterministically) initialize an integer array with values in [0, 9].
 - add a loop that sums even elements and subtracts odd elements.
 - visually check that it is correct.
 - Q: why is it not possible to initialize the array with any random integer value?
- Ex. 2: declare a synchronous channel and create two processes:
 - The first process sends the characters 'a' through 'z' onto the channel.
 - The second process reads the values of the channel and outputs them as characters.
 - Check if sooner or later the second process will read the letter 'z'.
- Ex. 3: replace the synchronous channel with a buffered channel and check how the behaviour changes.