

Grammars and Parsing

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1. Context Free Grammars (CFGs)

2. Efficiency and Expressivity

3. Features and Unification

4. Dependency Grammars

5. Resolving Ambiguity

6. Treebanks and Evaluation

1. Context Free Grammars

- Grammars
- Context Free Grammars (CFGs)
- Basic Parsing Strategies for CFGs
 - Top-Down
 - Bottom-Up
- Parsing and Search
- Redundancy in Parsing

Grammars

- A *grammar* is a 4-tuple $G = (N, \Sigma, P, S)$, where
 - N is a finite set of *nonterminal symbols*
 - Σ is a finite set of *terminal symbols*, disjoint from N
 - P is a set of rules, i.e. a finite subset of
$$(N \cup \Sigma)^* N (N \cup \Sigma)^* \times (N \cup \Sigma)^*$$Productions $(\alpha, \beta) \in P$ are usually written $\alpha \rightarrow \beta$
 - S is a distinguished symbol in N called the start symbol

Chomsky hierarchy

Different types of grammars/languages according to the definition of P :

- Regular grammars/languages
- Context-Free grammars/languages
- Context-Sensitive grammars/languages
- Unrestricted grammars/languages

Rules

- Regular Grammars:

$$A \rightarrow xB$$

$$A \rightarrow x$$

where A and B are in N and x is in Σ^*

- Context-Free Grammars:

$$A \rightarrow \alpha$$

where A is in N and α is in $(N \cup \Sigma)^*$

- Context-Sensitive Grammars:

$$\alpha \rightarrow \beta$$

where $|\alpha| \leq |\beta|$

Phrase Structure

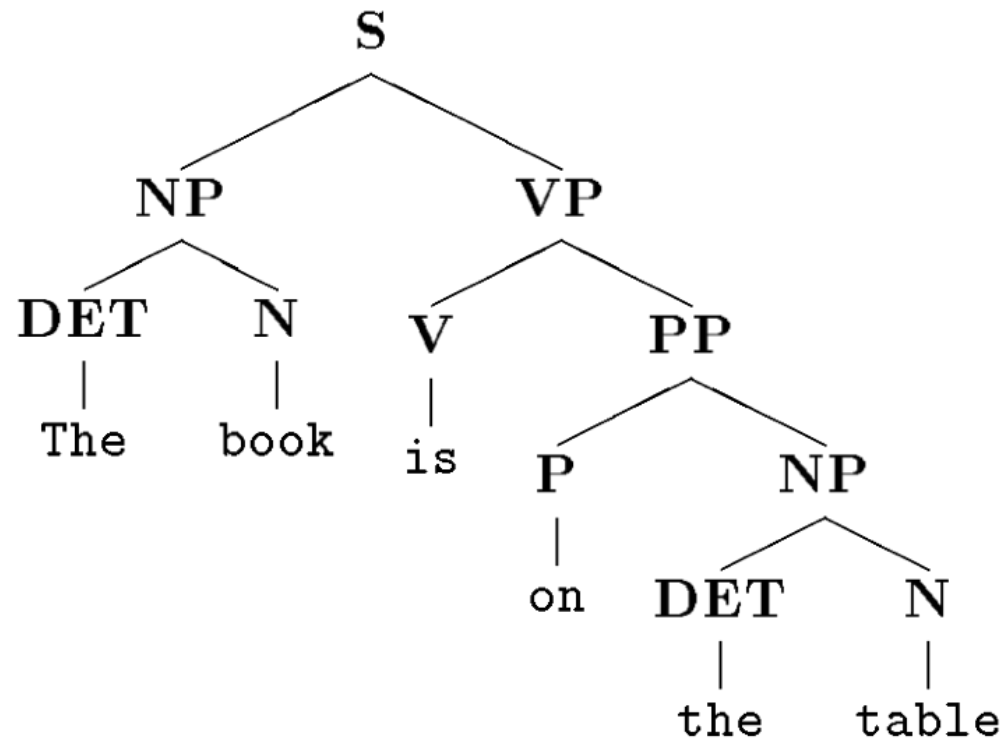
- *Language* = collection of *strings*
but ...
- Importance of hierarchical structure as well as linear structure of a given sentence

the book is on the table

Sentence:

the book is on the table

Parse tree:



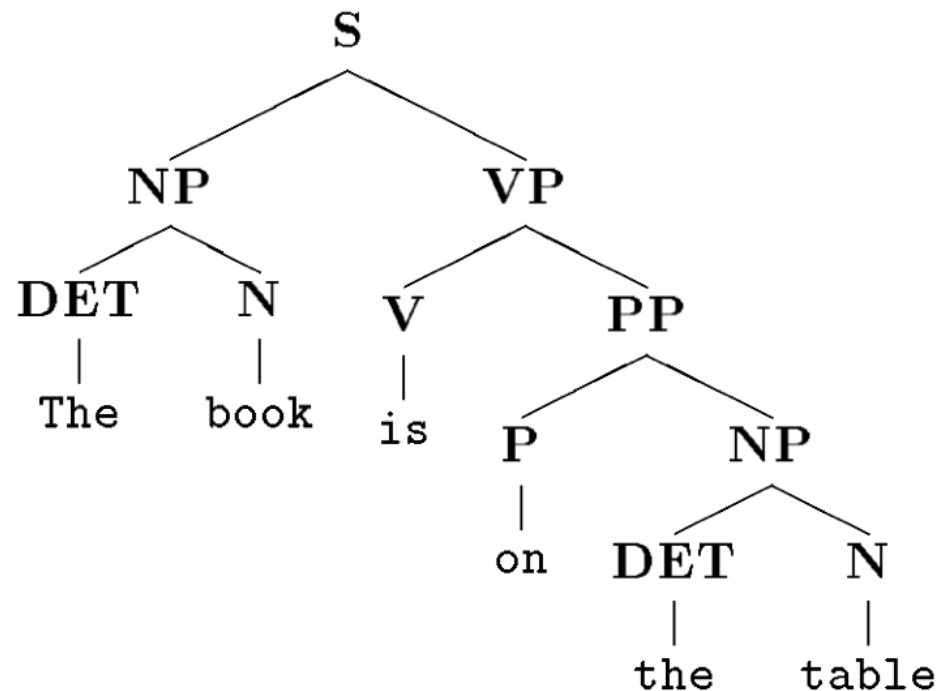
- Lexical elements:
 - *the* (DET)
 - *book, table* (Noun)
 - *is* (Verb)
 - *on* (Preposition)
- Constituent phrases:
 - *the book* (Noun Phrase)
 - *the table* (Noun Phrase)
 - *on the table* (Prepositional Phrase)
 - ...

Phrase Structure

Constituents can be indicated either by bracketing

[_S [_{NP} [_{DET} *the*] [_N *book*]] [_{VP} [_V *is*] [_{PP} *on* [_{NP} [_{DET} *the*] [_N *table*]]]]]

or by means of parse trees



Phrase Structure

- Hierarchical information about constituents (dominance)
- Linear precedence information
- Labelling information (syntactic categories)

Context-Free Grammars

Phrase structure grammars (PSGs) provide a means of characterizing the structure of sentences

A Context-Free (Phrase Structure) Grammar consists of a set of rules of the following form:

$$A \rightarrow X_1 X_2 \dots X_k \quad (k \geq 0)$$

- A is a *nonterminal* (a category name; e.g. N, NP, VP, DET, etc.)
- each X_i is either a *nonterminal* or a terminal (i.e. a word)

An example of a simple CFG

1. $S \rightarrow NP VP$

2. $NP \rightarrow John$

3. $NP \rightarrow Mary$

4. $NP \rightarrow DET N$

5. $DET \rightarrow a$

6. $N \rightarrow letter$

7. $VP \rightarrow V NP$

8. $VP \rightarrow V NP PP$

9. $VP \rightarrow V PP$

10. $V \rightarrow wrote$

11. $PP \rightarrow P NP$

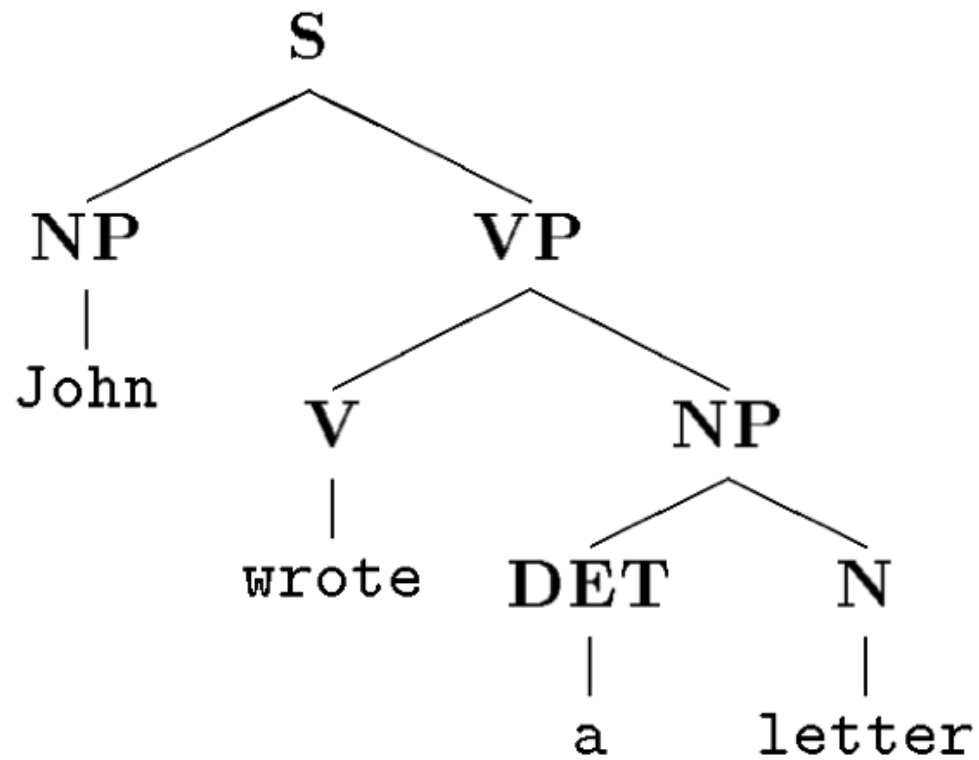
12. $P \rightarrow to$

John wrote a letter

John wrote a letter to Mary

John wrote to Mary

...



Three questions

- Are there effective procedures for recognition/generation of CFGs?
- How do we use CFGs to *parse* (i.e. assign structure to) strings?
- How do CFGs compare with FSLs computationally/descriptively?

Basic Parsing Strategies

Top-Down: *A goal-driven* strategy:

1. assume you are looking for S (i.e. sentence);
2. use rules ‘forward’ to ‘expand’ symbols until input is derived (else **fail**)

Bottom-Up: *A data-driven* strategy:

1. assume you are looking for S ;
2. use rules ‘backward’ to ‘combine’ symbols until you get S (else **fail**)

Basic Parsing Strategies

Other dimensions:

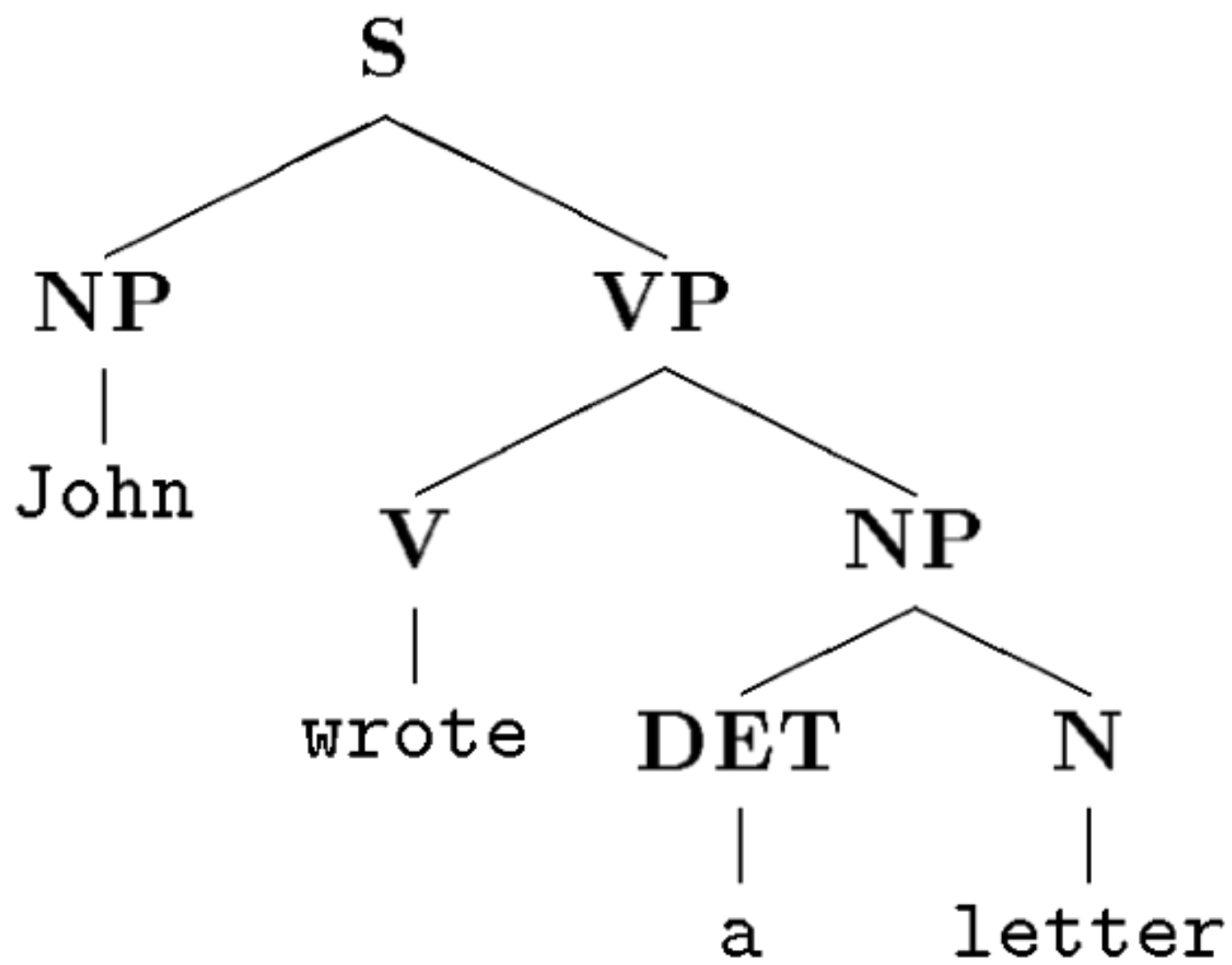
- *left-to-right* vs. *right-to-left* (but also *head-driven* or *island-driven*)
- *depth-first* vs. *breadth-first*

In the following examples, *left-to-right* and *depth-first* are usually adopted

Top-Down Strategy

Input: *John wrote a letter*

1	S	:	<i>John wrote a letter</i>	
2	NP VP	:	<i>John wrote a letter</i>	S \rightarrow NP VP
3	VP	:	<i>wrote a letter</i>	NP \rightarrow <i>John</i>
4	V NP	:	<i>wrote a letter</i>	VP \rightarrow V NP
5	NP	:	<i>a letter</i>	V \rightarrow <i>wrote</i>
6	DET N	:	<i>a letter</i>	NP \rightarrow DET N
7	N	:	<i>letter</i>	DET \rightarrow <i>a</i>
8	:			N \rightarrow <i>letter</i>



Crucial Points (1)

Non-determinism: at step 4, we could have chosen to ‘expand’ VP according to rule 8:

3 VP : *wrote a letter*

4 V NP PP : *wrote a letter* VP \rightarrow V NP PP

Need some way of exploring the possibilities and recovering if necessary (backtracking)

Crucial Points (2)

Left recursion: a problem for top-down strategy.

If we added a new rule:

$$(13) \text{ VP} \rightarrow \text{VP PP}$$

3	VP	:	<i>wrote a letter</i>	
4	VP PP	:	<i>wrote a letter</i>	VP \rightarrow VP PP
5	VP PP PP	:	<i>wrote a letter</i>	VP \rightarrow VP PP

and so on...

Bottom-Up Strategy

Input: *John wrote a letter*

1	<i>John wrote a letter</i>	
2	NP <i>wrote a letter</i>	NP \rightarrow <i>John</i>
3	NP V <i>a letter</i>	V \rightarrow <i>wrote</i>
4	NP V DET <i>letter</i>	DET \rightarrow <i>a</i>
5	NP V DET N	N \rightarrow <i>letter</i>
6	NP V NP	NP \rightarrow DET N
7	NP VP	VP \rightarrow V NP
8	S	S \rightarrow NP VP

Crucial Points

Empty productions: a problem for bottom-up strategy.
Empty productions have the form:

$$A \rightarrow \varepsilon$$

E.g.:

$NP \rightarrow DET AP N$

$AP \rightarrow \varepsilon$

$AP \rightarrow ADJ AP$

$ADJ \rightarrow \textit{lengthy}$

$ADJ \rightarrow \textit{interesting}$

Crucial Points

These new rules allow NPs such as:

a lengthy letter

a lengthy interesting letter

a letter

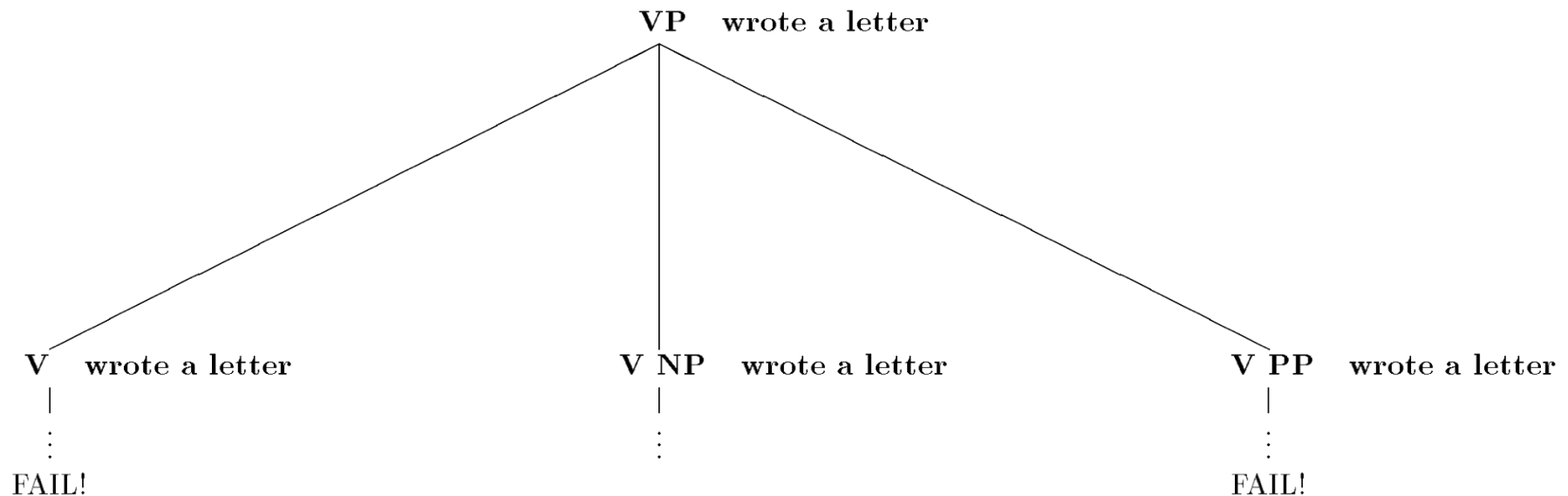
Note, however, that the rule $AP \rightarrow \varepsilon$ is always applicable!

Parsing and Search

In general, CFG parsing is *non-deterministic*.

Top-Down Example:

At different stages in the parsing process, more than one rule may be applicable:



Parsing and Search

Parsing algorithms need to explore the search space systematically.

To recover from errors, it is necessary to record the state of a parse each time a choice occurs.

Parsing and Search

E.g., considering the previous example,

The parse state:

VP : *wrote a letter*

has three different successor states:

V : *wrote a letter*

V NP : *wrote a letter*

V NP PP : *wrote a letter*

Parsing and Search

Parse maintains a list of parse states called an *agenda*:

- remove states from agenda;
- generate successor states;
- add successors to agenda;

Parse terminates successfully if the goal state (:) is generated.

Parse terminates unsuccessfully if it runs out of parse states to explore (i.e. the agenda is empty).

Parsing and Search

Search strategy:

this is determined by the order in which agenda items are considered:

Rule $S \rightarrow S_1 S_2 \dots S_k$

σ = rest on agenda

Depth-first:

$$S \sigma \Rightarrow S_1 S_2 \dots S_k \sigma$$

Breadth-first:

$$S \sigma \Rightarrow \sigma S_1 S_2 \dots S_k$$

Redundancy in Parsing

Input: *John sang a song*

1	S	:	<i>John sang a song</i>	
2	NP VP	:	<i>John sang a song</i>	S → NP VP
3	VP	:	<i>sang a song</i>	NP → <i>John</i>
4	V NP PP	:	<i>sang a song</i>	VP → V NP PP
5	NP PP	:	<i>a song</i>	V → <i>sang</i>
6	DET N PP	:	<i>a song</i>	NP → DET N
7	N PP	:	<i>song</i>	DET → <i>a</i>
8	PP	:		N → <i>song</i>
9	P NP	:		PP → P NP

Redundancy in Parsing

4'	V NP	:	<i>sang a song</i>	VP → V NP
5'	NP	:	<i>a song</i>	V → <i>sang</i>
6'	DET N	:	<i>a song</i>	NP → DET N
7'	N	:	<i>song</i>	DET → <i>a</i>
8'		:		N → <i>song</i>

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2. Efficiency and Expressivity

- Efficiency
 - Redundancy in Parsing
 - Active Chart Parsing
- Expressivity
 - comparing CFGs and FSAs
 - Pros and Cons of CFGs
 - Agreement, subcategorization, ...

Chart Parsing

Dynamic programming technique which keeps track of what has been done and of partial hypotheses.

Resulting data structure is called the *active chart*.

The chart contains data structures called *edges*, which represent (partially) recognized constituents.

Dotted Rules

‘Dotted Rules’: edges have labels of the general form:

$$C \rightarrow X_1 \dots X_j \bullet X_{j+1} \dots X_k$$

Symbols on the left of the dot (\bullet) have been already ‘found’ (confirmed hypotheses). Symbols on the right are still to be found.

Chart Parsing

NP \rightarrow DET • N

active edge

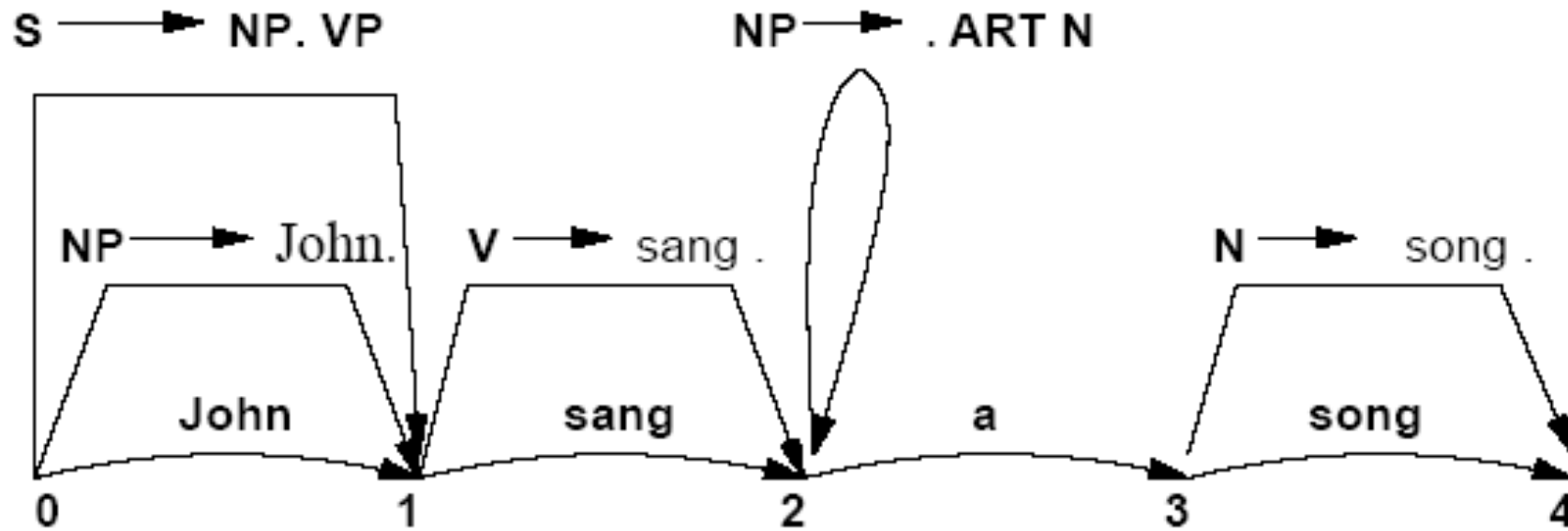
S \rightarrow • NP VP

active and empty edge

VP \rightarrow V NP •

inactive edge

Chart Parsing



The chart has edges of the form

$$(i, j, A \rightarrow \alpha \bullet \beta)$$

Fundamental Rule of Chart Parsing

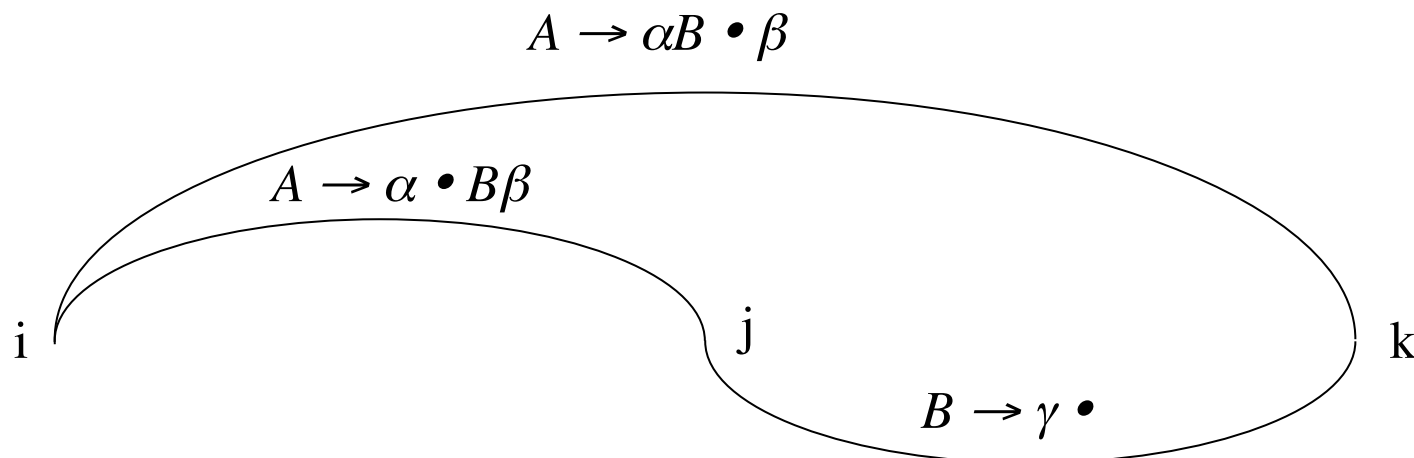
IF the chart contains the edges:

$$(i, j, A \rightarrow \alpha \bullet B\beta) \text{ and } (j, k, B \rightarrow \gamma \bullet)$$

THEN add the new edge:

$$(i, k, A \rightarrow \alpha B \bullet \beta)$$

(α , β , γ possibly empty strings of symbols)



Fundamental Rule of Chart Parsing

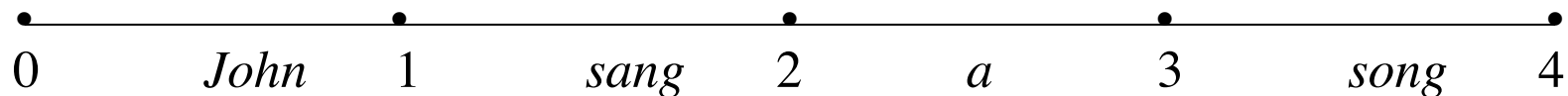
Fundamental rule only applies to chart containing active and inactive edges.

- How do we get started

Initialization:

Initially chart contains inactive edges corresponding to words in the input string:

e.g. for input *John sang a song*



Rule Invocation

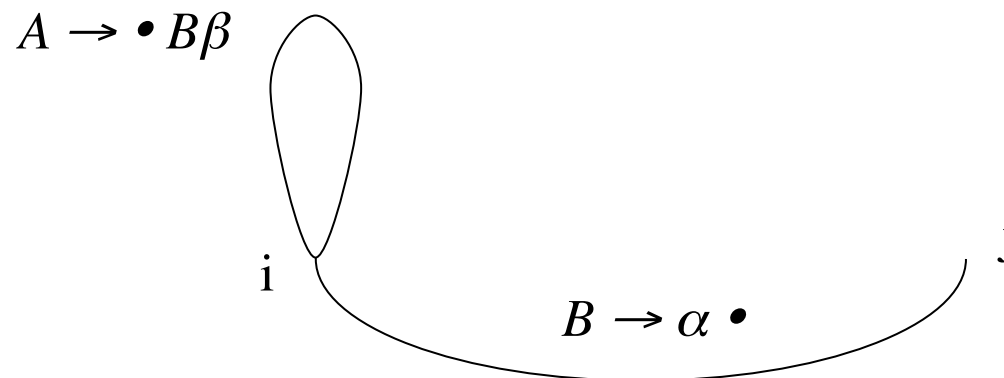
Bottom-Up :

IF you add an edge $(i, j, B \rightarrow \alpha \bullet)$

THEN for every rule of the form

$$A \rightarrow \bullet B\beta$$

add an edge $(i, i, A \rightarrow \bullet B\beta)$



Rule Invocation

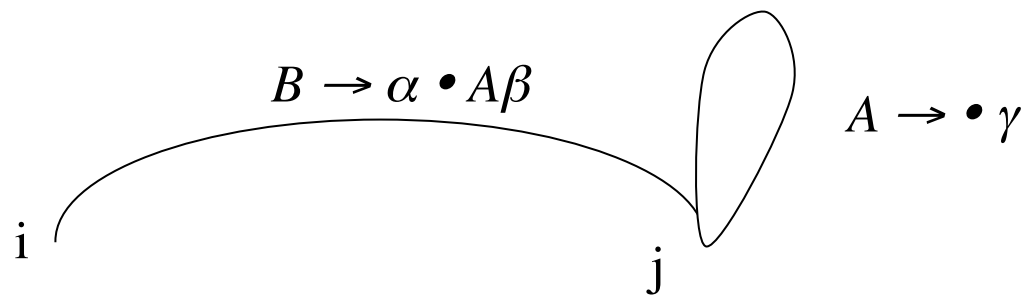
Top-Down:

IF you add an edge $(i, j, B \rightarrow \alpha \bullet A\beta)$

THEN for every rule of the form

$$A \rightarrow \gamma$$

add an edge $(j, j, A \rightarrow \bullet \gamma)$



Comparing CFGs and FSAs

FSAs:

- recognition is efficient – linear time; but
- the formalism is not very expressive.

CFGs:

- the basic parsing (recognition) strategies are not efficient – exponential time; but
- using dynamic programming techniques we can do better than this (Chart parsing; CKY algorithm; Earley algorithm); and
- CFGs are more expressive than FSAs

Comparing CFGs and FSAs

- Any language describable with a FSA is describable with a CFG.
- There are languages that can be described with a CFG that cannot be described with a FSA.

Regular \subset Context Free

- There is a general agreement that NLs are not Regular languages (i.e. cannot be adequately described with FSAs)
- Much of the syntax of the world's NLs seems to be Context Free (i.e. can be adequately described with CFGs).

Pros and Cons of CFGs

Advantages:

- Can describe infinite languages and assign appropriate syntactic structures
- Recognition (parsing) procedures can be implemented reasonably efficiently – $O(n^3)$:
 - Earley algorithm (Chart Parsing)
 - Cocke-Kasami-Younger (CKY) algorithm
 - Tomita's generalized LR parser

Pros and Cons of CFGs

- NLS \cong CFGLs?
 - Long-standing argument
 - Arguably some NLS are non-CFLs (e.g. Swiss German – Shieber 1985)

Pros and Cons of CFGs

Disadvantages:

- Difficult to capture certain NL phenomena appropriately/adequately/elegantly:
 - agreement
 - subcategorization
 - generalizations over word/constituent order
 - relationships between different sentence types
- Some NL phenomena appear to require greater mathematical expressivity (i.e. there is evidence that some NLs are not CFLs)

Grammar equivalence

- Two grammars are **weakly equivalent** if they generate the same language (i.e. the same set of strings)
- Two grammars are **strongly equivalent** if they generate the same language *and* they assign the same phrase structure to each sentence

Mildly context-sensitive grammars (e.g. TAGs, Tree Adjoining Grammars)

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3. Features and Unification

- Limitations of CFGs
- Unification-Based Grammars
 - Feature Structures
 - Unification
 - The PATR Formalism
 - Typed Feature Structures

Agreement phenomena

Verbs have to “agree” with subjects

NP	VP	
<i>the boy</i>	<i>sees the girl(s)</i>	singular
<i>the boys</i>	<i>see the girl(s)</i>	plural

Agreement phenomena

$S \rightarrow NP_s VP_s$

$S \rightarrow NP_p VP_p$

$NP_s \rightarrow DET_s N_s$

$NP_p \rightarrow DET_p N_p$

$DET_s \rightarrow the$

$DET_p \rightarrow the$

$N_s \rightarrow boy$

$N_p \rightarrow boys$

$VP_s \rightarrow V_s NP_s$

$VP_s \rightarrow V_s NP_p$

$VP_p \rightarrow V_p NP_s$

$VP_p \rightarrow V_p NP_p$

$V_s \rightarrow sees$

$V_p \rightarrow see$

$N_s \rightarrow girl$

$N_p \rightarrow girls$

Subcategorization

Different verbs may require different complements

VP → V1	(<i>die</i>)
VP → V2 NP	(<i>love</i>)
VP → V3 NP NP	(<i>give</i>)
VP → V4 NP PP	(<i>put</i>)
VP → V5 NP S	(<i>tell</i>)
VP → V6 S	(<i>believe</i>)
and so on...	

Unbounded Dependency Constructions

E.g. Wh-questions:

Who did Bill see ϵ ?

Who did Tom say that Bill saw ϵ ?

Who did Anna believe Tom said that Bill saw ϵ ?

- Correct interpretation of *who* depends on structure which is arbitrarily distant
- Difficult to capture UDCs with simple CFGs

Criteria for Formalism Design in NLP

Generative Power: can the formalism describe the language at all?

Notational Expressivity: can the formalism capture the appropriate generalizations?

Computational Effectiveness: does the formalism have a sensible, practical computational interpretation?

Note: simple CFGs score quite well on the first and third criteria; less well on the second.

Unification-Based Grammars

A family of related grammar formalisms.

UBGs can be viewed as extensions of CFGs which

- make use of constraints on feature values (to capture agreement, etc.)
- make use of syntactic features and allow underspecification of linguistic objects (categories or other representations)
- employ unification as a consistency checking / information merging operation

Examples of UBGs

- FUG (Kay)
- LFG (Bresnan & Kaplan)
- GPSG (Gazdar, Klein, Pullum & Sag)
- HPSG (Pollard & Sag)
- PATR (Shieber)
- CUG (Uszkoreit)
- UCG (Calder et al.)
- DUG (Hellwig)
- RUG (Carlson)
- TUG (Popovich)

Feature Structures

UBGs employ record-like objects to represent categories.

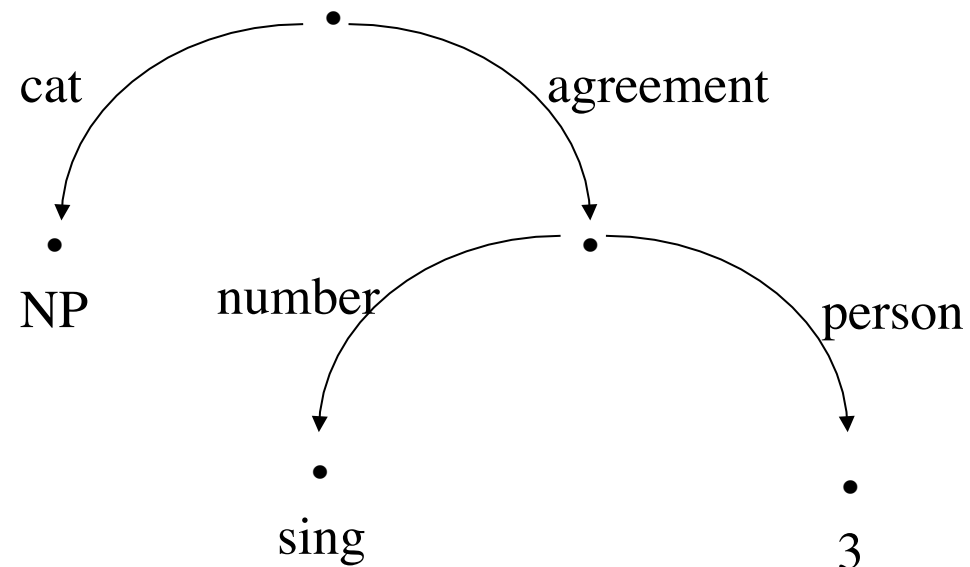
Third person singular NP:
$$\left[\begin{array}{ll} \text{CAT} & \text{NP} \\ \text{AGREEMENT} & \left[\begin{array}{ll} \text{NUMBER} & \text{sing} \\ \text{PERSON} & 3 \end{array} \right] \end{array} \right]$$

- made up of features (cat, agreement, number, person) and values
- values may be *simple* (e.g., NP, sing and 3) or *complex*:

$$\left[\begin{array}{ll} \text{NUMBER} & \text{sing} \\ \text{PERSON} & 3 \end{array} \right]$$

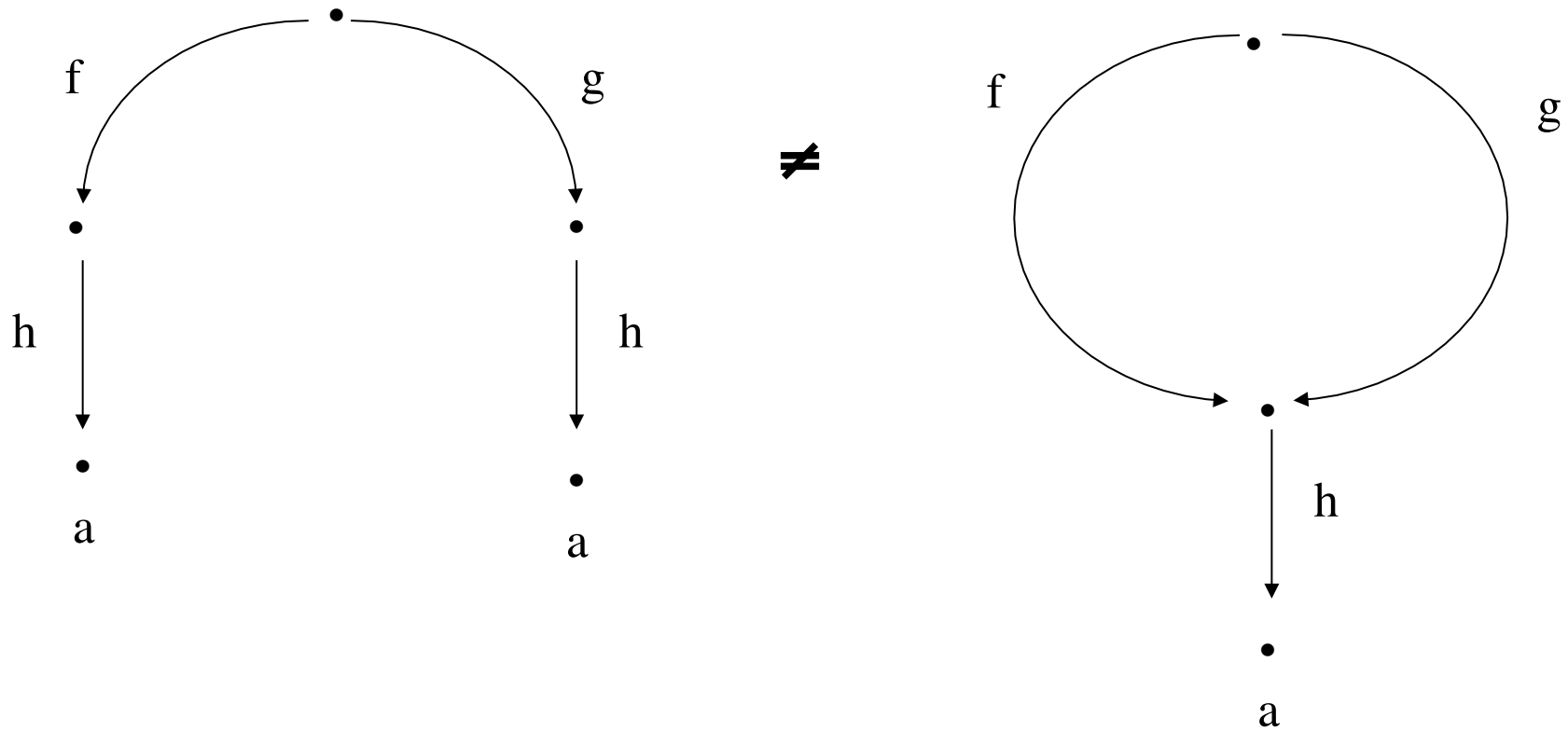
Feature Structures

Feature structures may be drawn as directed graphs:



Feature Structures

Feature structures may be *re-entrant*:



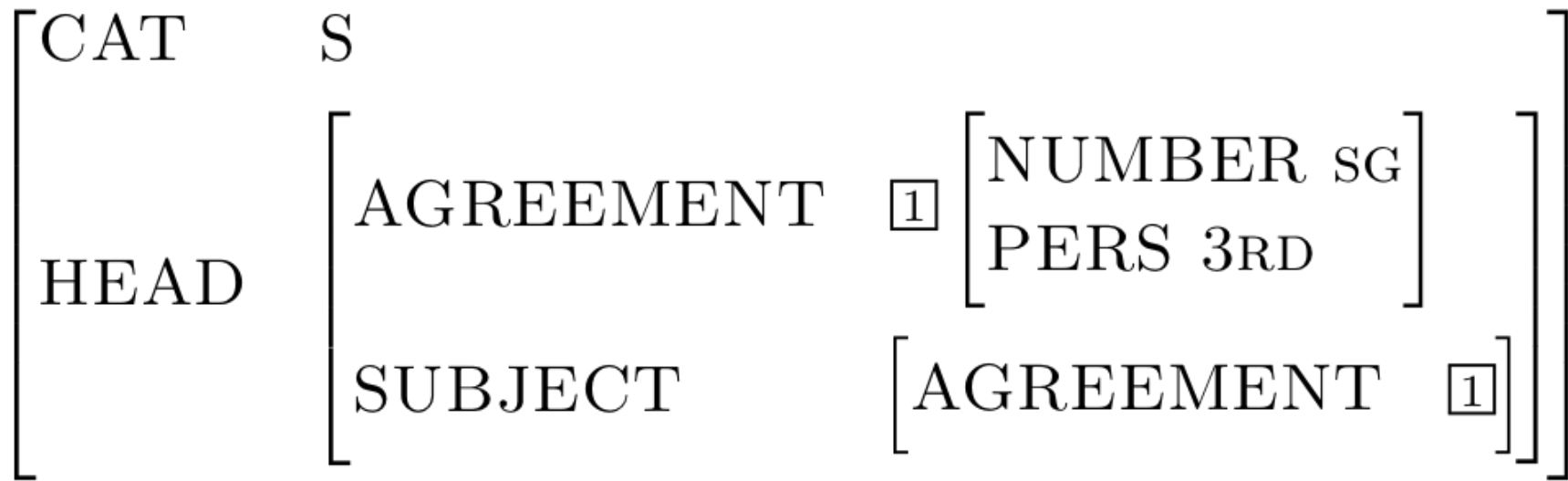
Feature Structures

Feature structures may be *re-entrant*:

$$\begin{bmatrix} \text{F} & \begin{bmatrix} \text{H} & \text{a} \end{bmatrix} \\ \text{G} & \begin{bmatrix} \text{H} & \text{a} \end{bmatrix} \end{bmatrix} \neq \begin{bmatrix} \text{F} & \boxed{1} & \begin{bmatrix} \text{H} & \text{a} \end{bmatrix} \\ \text{G} & \boxed{1} & \end{bmatrix}$$

Reentrant Feature Structures

A linguistic example:



Feature Structures

Feature structures allow for underspecification of categories

Singular NP:

$$\left[\begin{array}{ll} \text{CAT} & \text{NP} \\ \text{AGREEMENT} & \left[\text{NUMBER} \quad \text{sing} \right] \end{array} \right]$$

Nominative NP:

$$\left[\begin{array}{ll} \text{CAT} & \text{NP} \\ \text{CASE} & \text{nom} \end{array} \right]$$

Unification

F1 =

$$\left[\begin{array}{l} \text{CAT} \quad \text{NP} \\ \text{AGREEMENT} \quad \left[\text{NUMBER} \quad \text{sing} \right] \end{array} \right]$$

F2 =

$$\left[\begin{array}{l} \text{CAT} \quad \text{NP} \\ \text{AGREEMENT} \quad \left[\text{PERSON} \quad 3 \right] \\ \text{CASE} \quad \text{nominative} \end{array} \right]$$

F1 \cup F2 =

$$\left[\begin{array}{l} \text{CAT} \quad \text{NP} \\ \text{AGREEMENT} \quad \left[\begin{array}{l} \text{NUMBER} \quad \text{sing} \\ \text{PERSON} \quad 3 \end{array} \right] \\ \text{CASE} \quad \text{nominative} \end{array} \right]$$

Unification

Unification fails when feature structures are incompatible

F3 =

CAT	VP				
AGREEMENT	<table><tr><td>NUMBER</td><td>sing</td></tr><tr><td>PERSON</td><td>3</td></tr></table>	NUMBER	sing	PERSON	3
NUMBER	sing				
PERSON	3				

F4 =

CAT	VP		
AGREEMENT	<table><tr><td>NUMBER</td><td>plu</td></tr></table>	NUMBER	plu
NUMBER	plu		
VFORM	tensed		

F3 \cup F4 = FAIL

The PATR Formalism

Originally introduced by Shieber and his colleagues at
SRI International

$$S \rightarrow NP VP$$
$$C_0 \rightarrow C_1 C_2$$
$$\langle C_0 \text{ cat} \rangle = S$$
$$\langle C_1 \text{ cat} \rangle = NP$$
$$\langle C_2 \text{ cat} \rangle = VP$$
$$\langle C_1 \text{ case} \rangle = \text{nominative}$$
$$\langle C_1 \text{ agreement} \rangle = \langle C_2 \text{ agreement} \rangle$$

Typed Feature Structures

Limitations of simple feature structure formalisms:

- No way to constrain possible values of a feature (e.g., the feature NUMBER can take only SING and PLU values)
- No way to capture generalization across feature structures (e.g., different types of English verb phrases)

Solution: use of **types**.

Typed Feature Structures

- Each feature structure is labeled by a type
- Each type has **appropriateness conditions** expressing which features are appropriate for it
- Types are organized in a **type hierarchy**
- Unification should take into account the types of feature structures in addition to unifying attributes and values

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4. Dependency Grammars

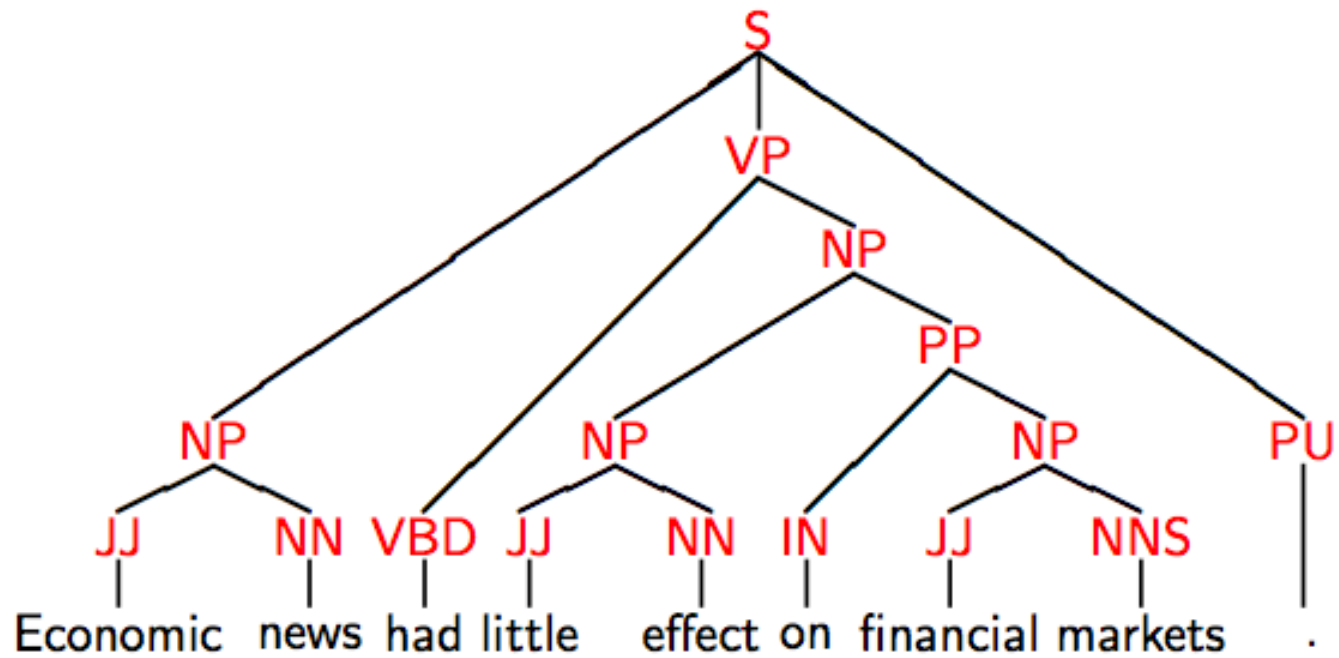
- Constituency vs. Dependency
- Dependency Grammars
- Dependency Parsing

some material taken from McDonald & Nivre ESLLI 2007 course
“Introduction on Data-Driven Dependency Parsing”

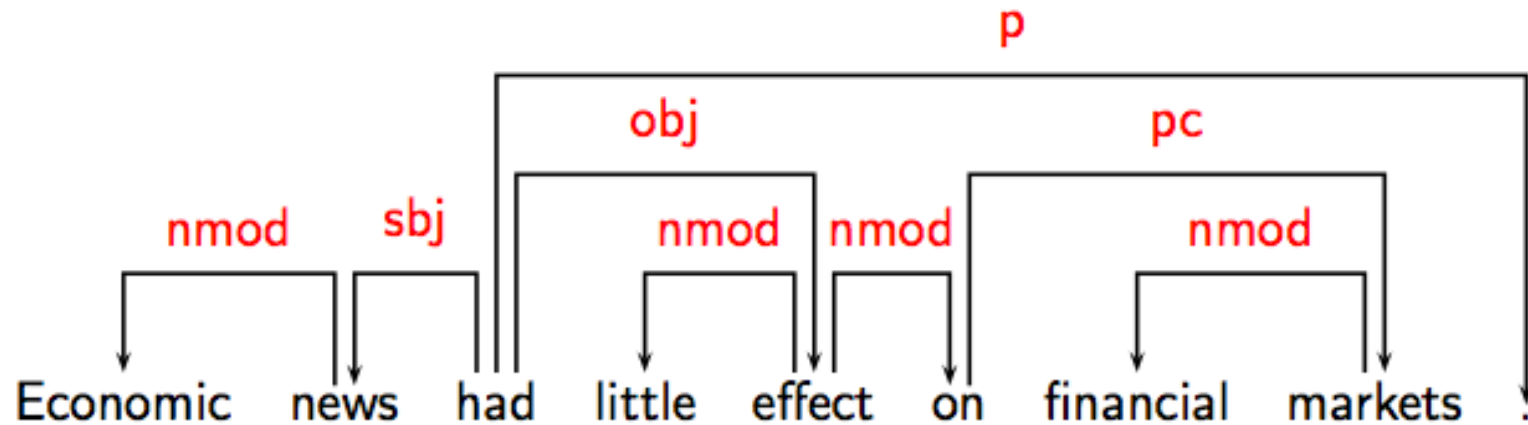
Dependency Grammars

- [Tesnière 1959]
- Syntactic structure of a sentence consists of **lexical items**, linked by binary asymmetric relations called **dependencies**.
- Dependency relations hold between a *head* (parent) and a *dependent* (daughter)

Phrase Structure



Dependency Structure



Constituency vs. Dependency

Phrase structure grammars

- Words appear only as leaves
- Internal nodes of trees consist of non-terminals

Dependency grammars

- No non-terminals
- Only words and binary relations between them

Constituency vs. Dependency

- Phrase structures explicitly represent
 - phrases (non-terminal nodes),
 - structural categories (non-terminal labels),
 - possibly some functional categories (grammatical functions).
- Dependency structures explicitly represent
 - head-dependent relations (directed arcs),
 - functional categories (arc labels),
 - possibly some structural categories (PoS).

Dependency Grammars

Family of grammatical formalisms differing in:

- Terminology (head/dependent, governor/modifier, regent/subordinate, ...)
- Criteria adopted to establish dependency relations
- Criteria to identify heads and dependents

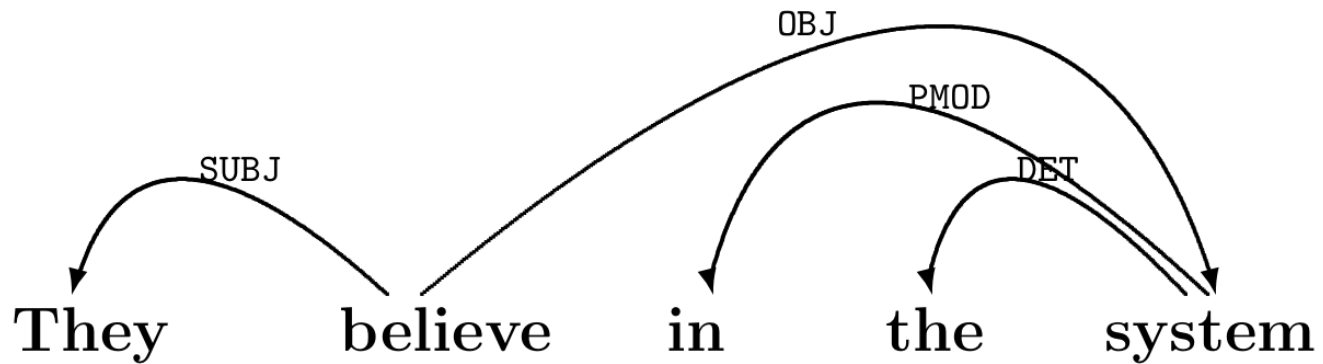
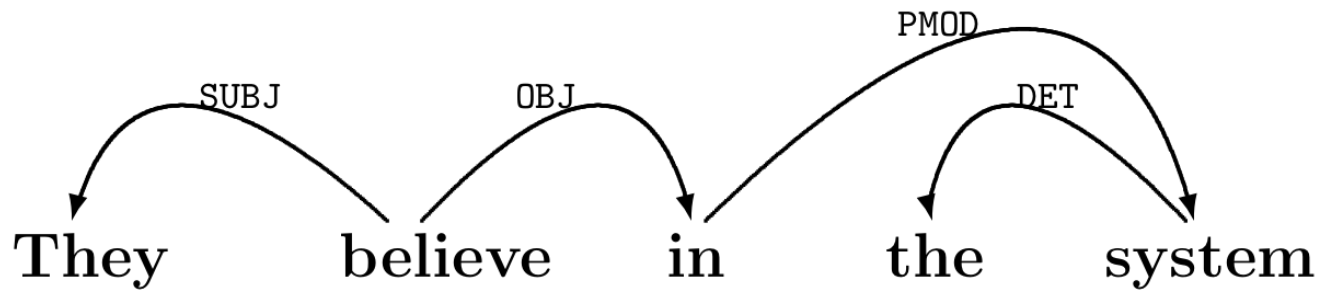
Dependency Relations

- Surface-oriented grammatical functions: *subject, object, adverbial, ...*
- Semantically oriented roles: *agent, patient, goal, ...*

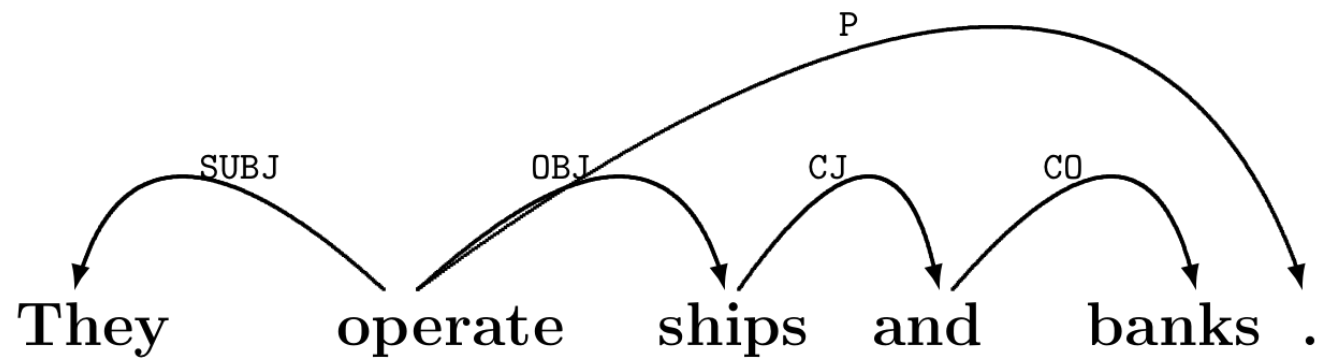
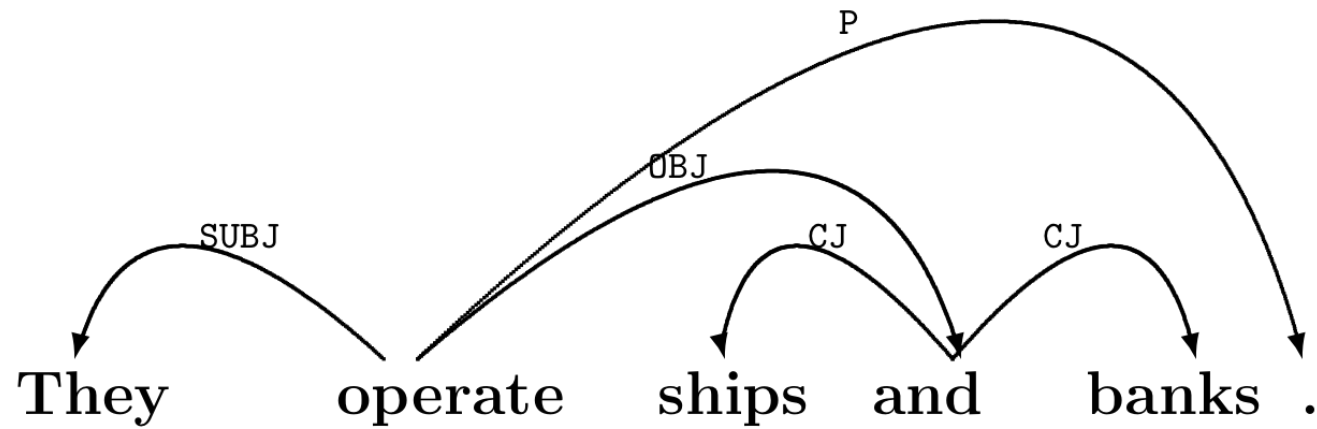
Problematic Constructions

- Grammatical function words: *syntactic* versus *semantic* heads
- Coordination: problematic in general

Function words



Coordination



Dependency Graphs

- A dependency structure can be defined as a directed graph G , consisting of
 - a set V of nodes (vertices),
 - a set A of arcs (directed edges),
 - a linear precedence order $<$ on V (word order).

Dependency Graphs

- Labeled graphs:
 - Nodes in V are labeled with word forms (and annotation).
 - Arcs in A are labeled with dependency types:
 - $L = \{l_1, \dots, l_{|L|}\}$ is the set of permissible arc labels.
 - Every arc in A is a triple (i, j, k) , representing a dependency from w_i to w_j with label l_k .

Dependency Parsing

- The problem:
 - Input: Sentence $x = w_0, w_1, \dots, w_n$ with $w_0 = \text{root}$
 - Output: Dependency graph $G = (V, A)$ for x where:
 - $V = \{0, 1, \dots, n\}$ is the vertex set,
 - A is the arc set, i.e., $(i, j, k) \in A$ represents a dependency from w_i to w_j with label $l_k \in L$